

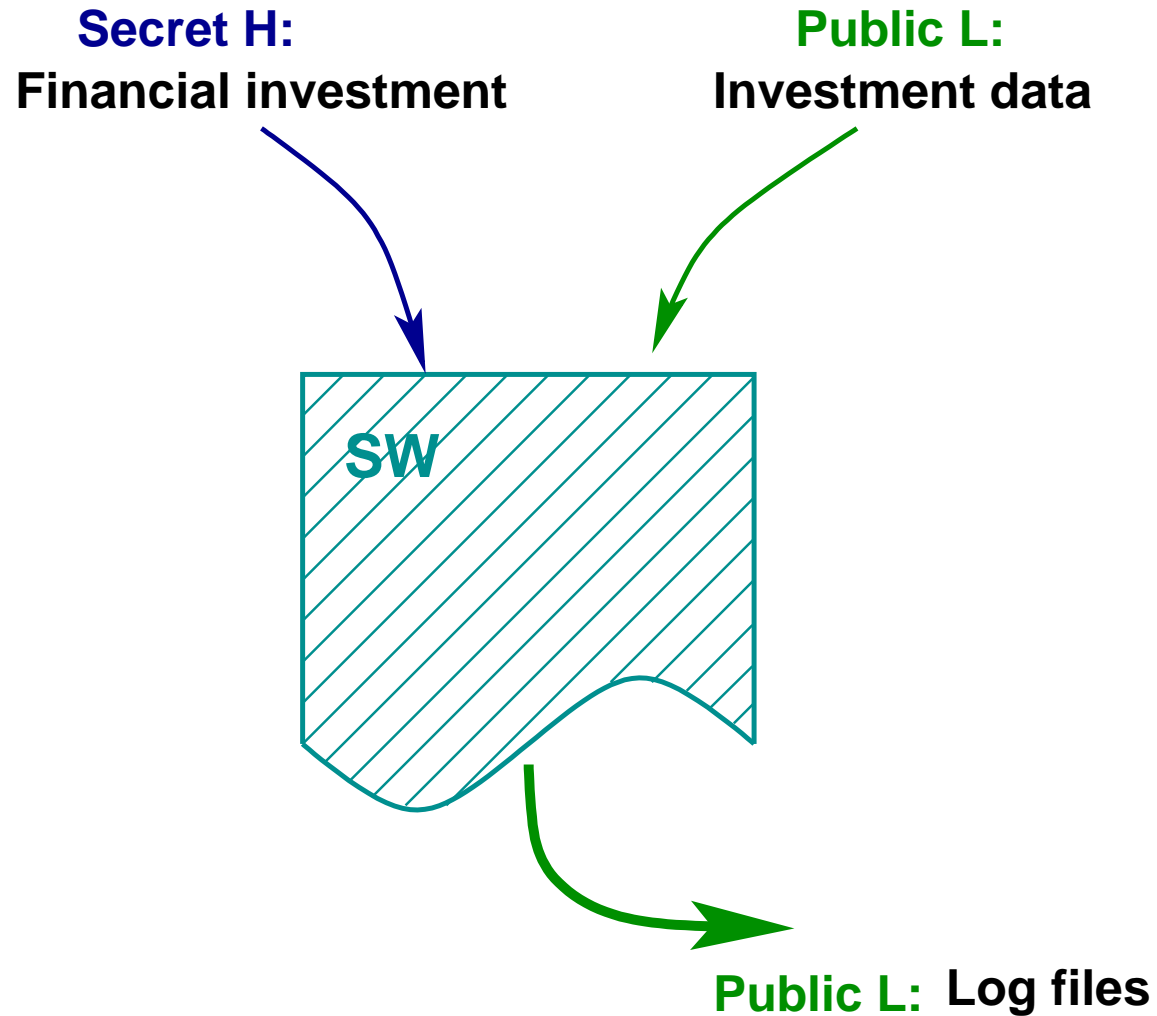
TIMED ABSTRACT NON-INTERFERENCE

Roberto Giacobazzi and Isabella Mastroeni

**Dipartimento di informatica
Università di Verona, Italy**

FORMATS'05, September 28, 2005

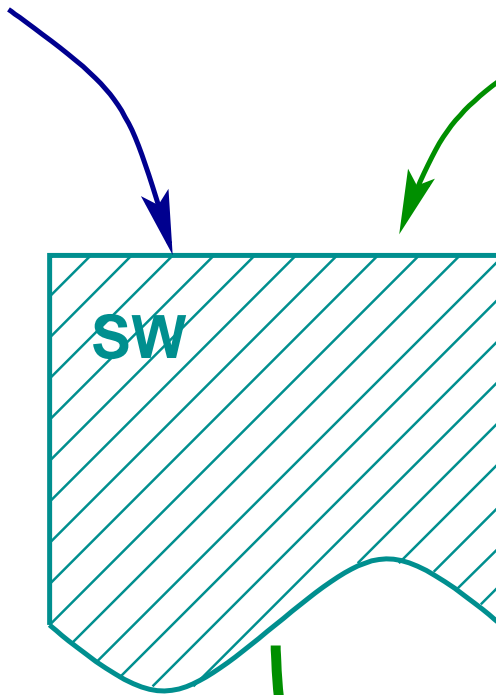
The Problem: Non-Interference



The Problem: Non-Interference

Secret H:
Financial investment

Public L:
Investment data



Is it secure?

Public L: Log files

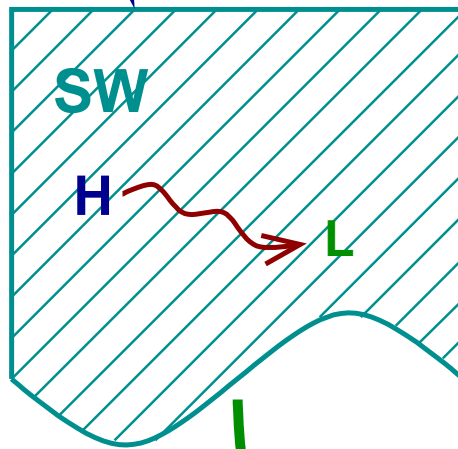


External observer

The Problem: Non-Interference

Secret H:
Financial investment

Public L:
Investment data



Is it secure? **NO**

Secret H
Public L: Log files



External observer

Background

SECURITY PROPERTY: States which classes have not to interfere with other classes of objects.

Background

SECURITY PROPERTY: States which classes have not to interfere with other classes of objects.



Confinement problem [Lampson'73]: *Preventing the results of computations leaking even partial information about the confidential inputs.*

Background

SECURITY PROPERTY: States which classes have not to interfere with other classes of objects.



Confinement problem [Lampson'73]: *Preventing the results of computations leaking even partial information about the confidential inputs.*



Non-interference policies require that any change upon confidential data has not to be revealed through the observation of public data.

- ⑥ Many real systems are intended to leak some kind of information
- ⑥ Even if a system satisfies non-interference, some kind of tests could reject it as insecure

Background

SECURITY PROPERTY: States which classes have not to interfere with other classes of objects.



Confinement problem [Lampson'73]: *Preventing the results of computations leaking even partial information about the confidential inputs.*

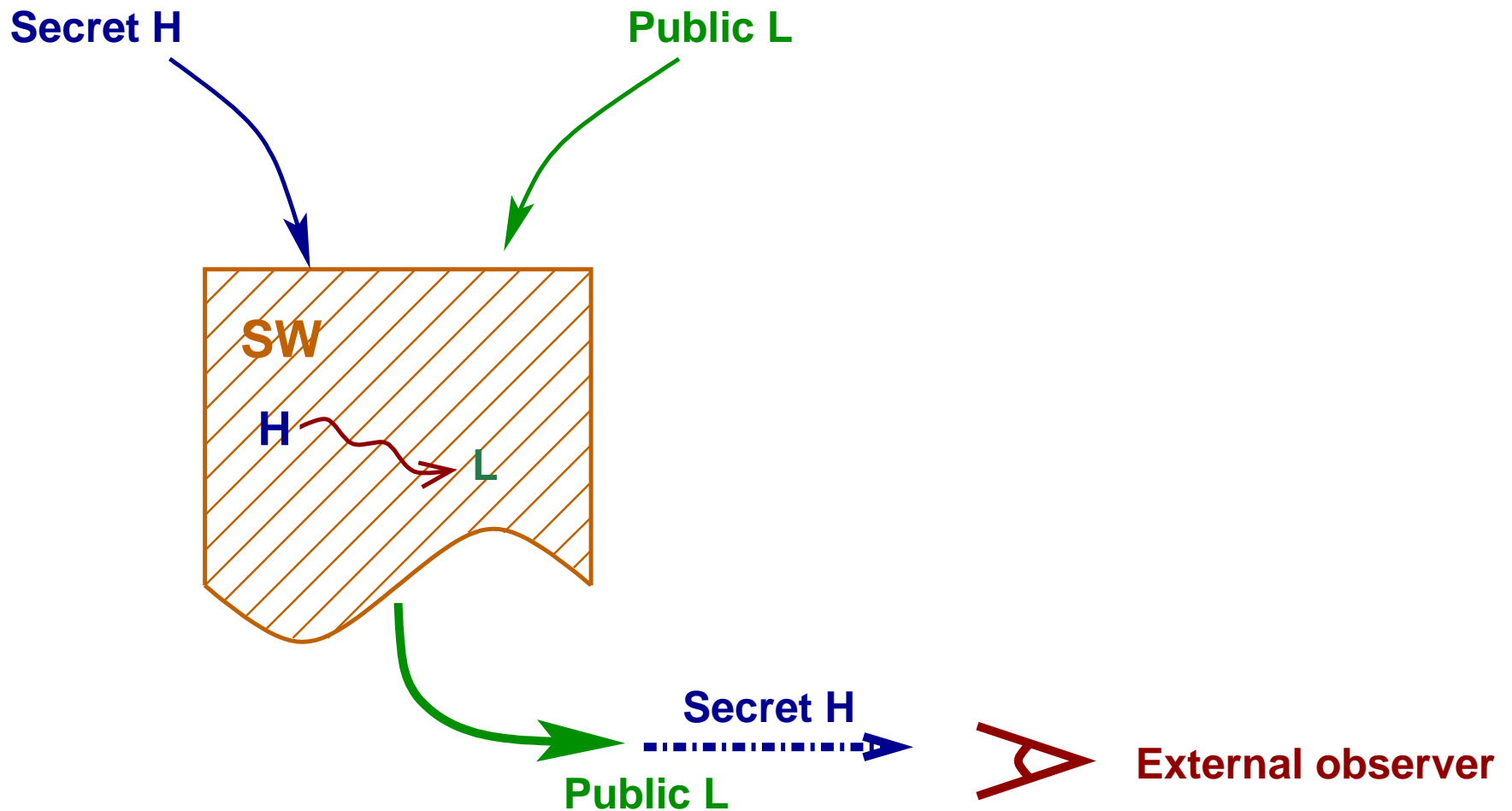


Non-interference policies require that any change upon confidential data has not to be revealed through the observation of public data.

- ⑥ **Characterizing released information:** [Cohen'77], [Zdancewic & Myers'01], [Clark et al.'04], [Lowe'02];
- ⑥ **Constraining attackers:** [Di Pierro et al.'02], [Laud'01].

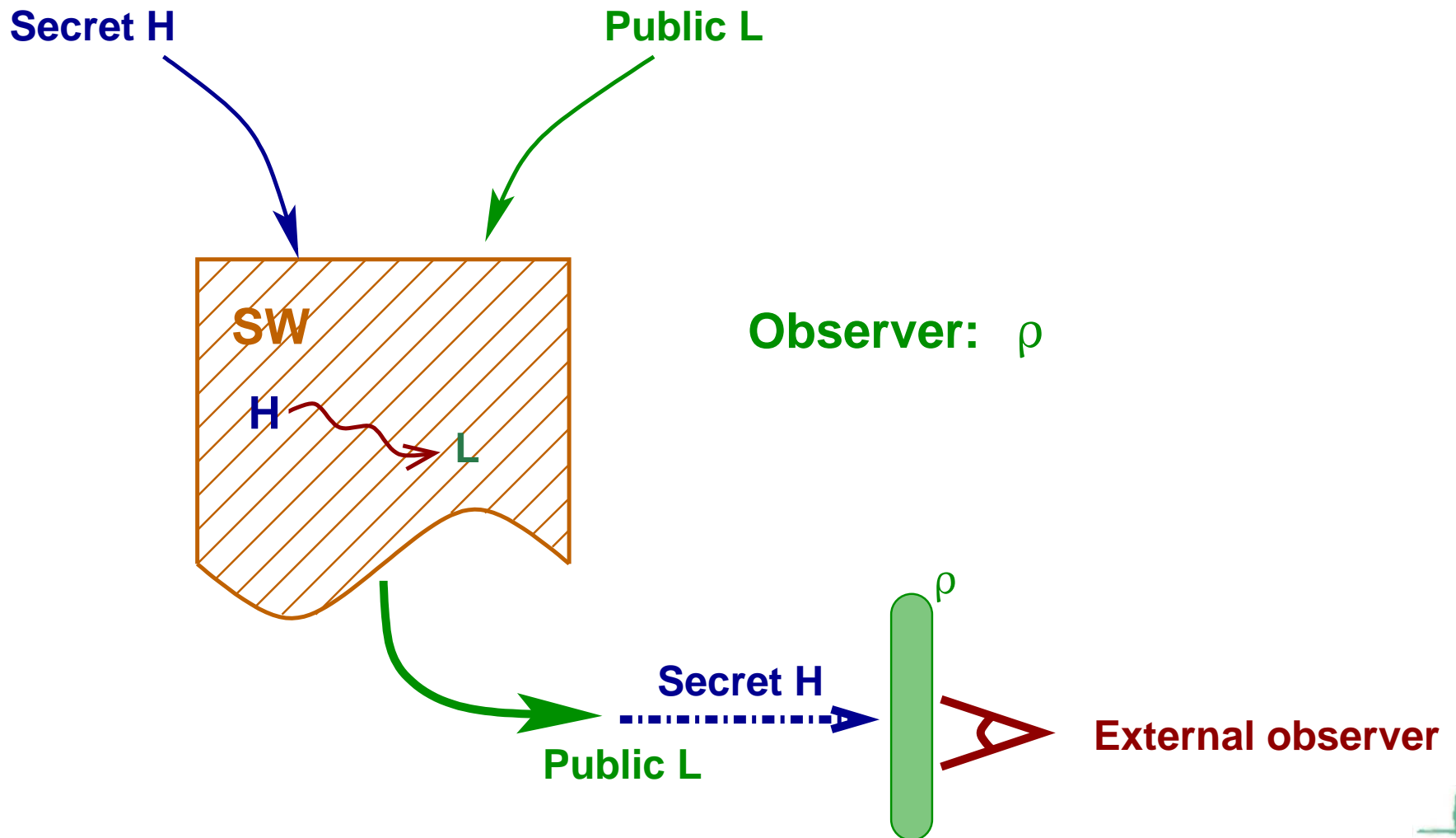
Abstracting Non-Interference

[Giacobazzi & Mastroeni, POPL'04]



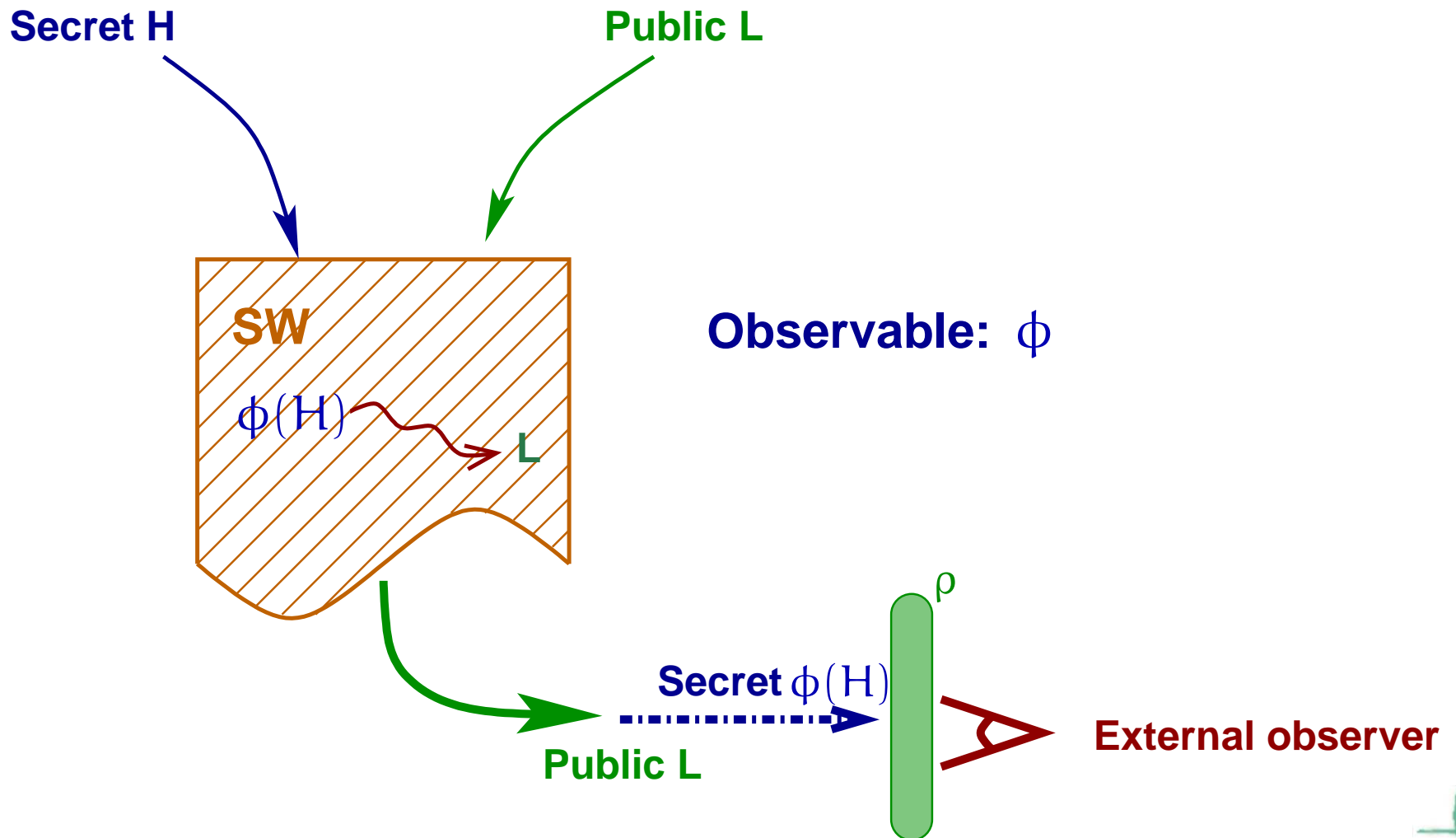
Abstracting Non-Interference

[Giacobazzi & Mastroeni, POPL'04]



Abstracting Non-Interference

[Giacobazzi & Mastroeni, POPL'04]



AI: Lattice of Abstractions

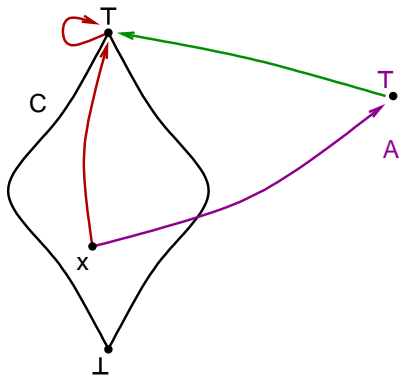
The concrete domain $\langle C, \leq, \wedge, \vee, \perp, \top \rangle$

[Cousot & Cousot '79]

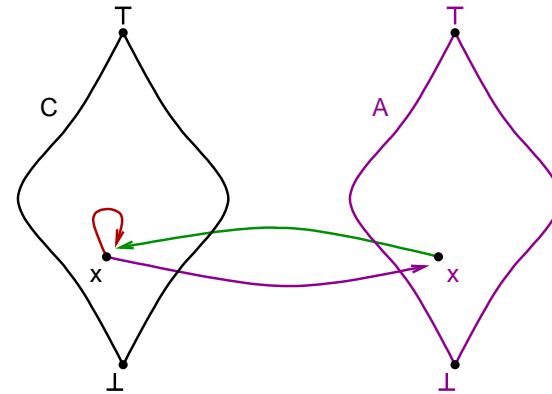
Lattice of abstract domains $\equiv \text{Abs}(C)$
 $\langle \text{Abs}(C), \sqsubseteq, \sqcap, \sqcup, \top, \perp \rangle$

$A_1 \sqsubseteq A_2 \Leftrightarrow A_2 \subseteq A_1$ (A_1 more precise than A_2)

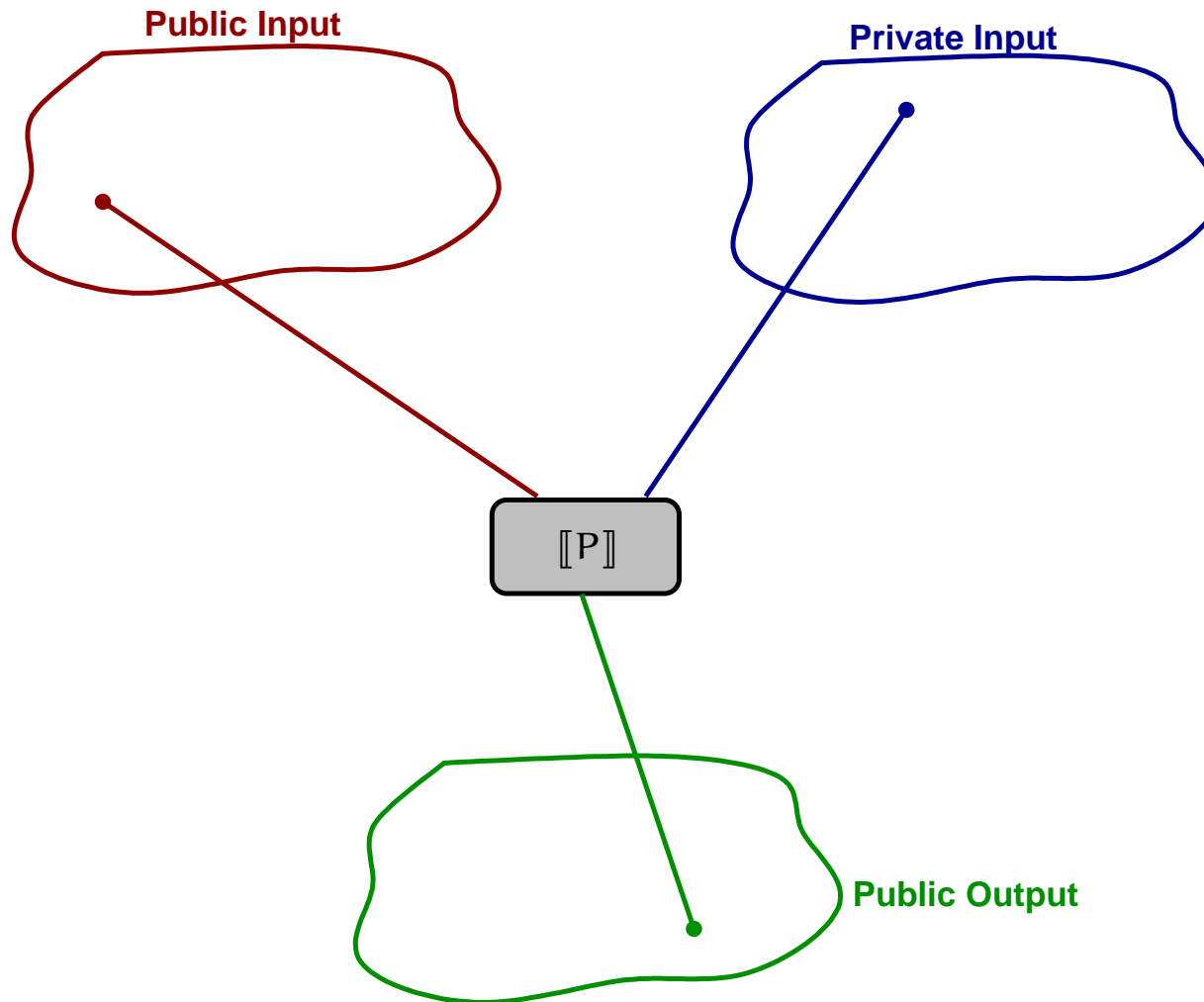
Top:



Bottom:

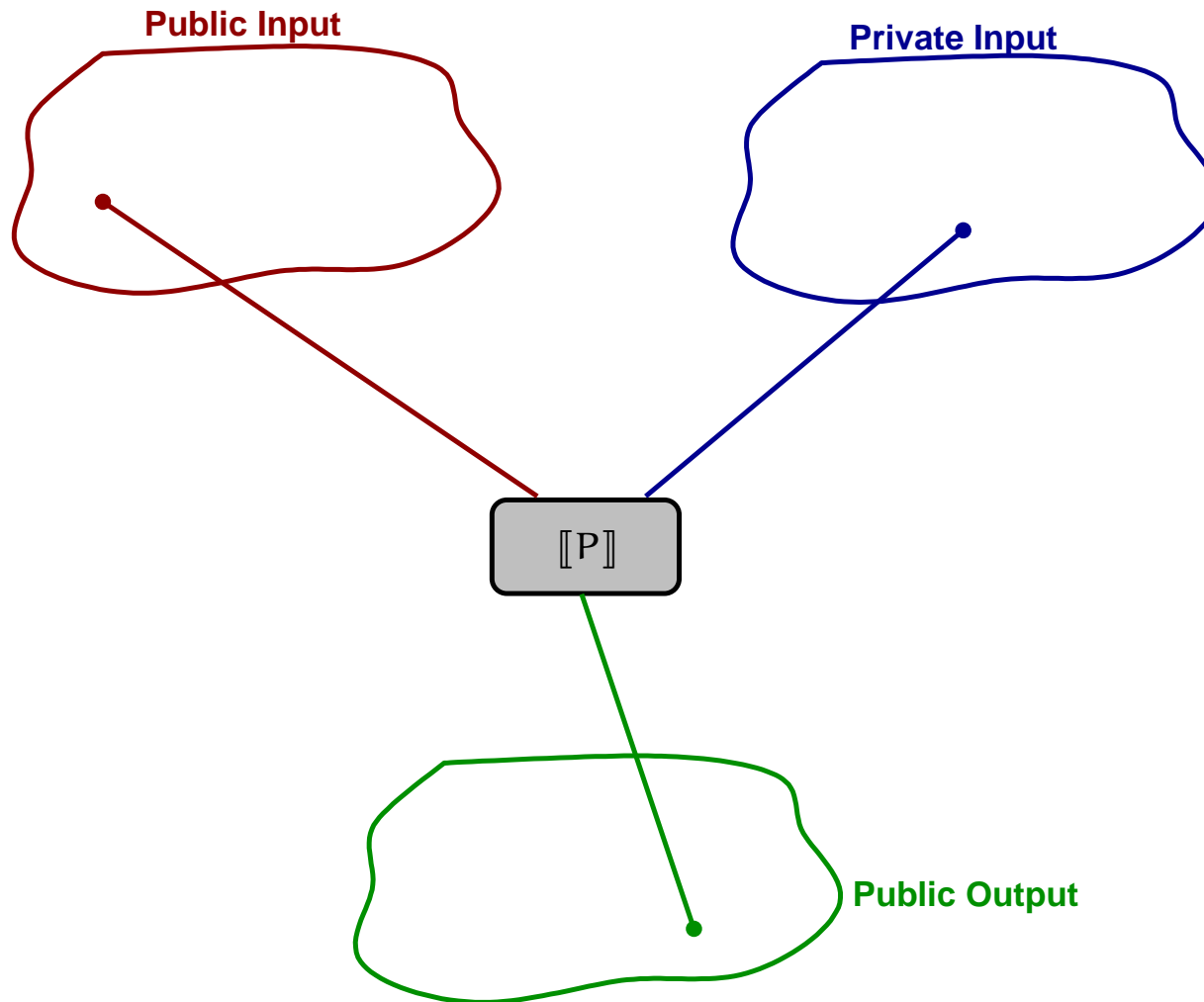


Standard non-interference



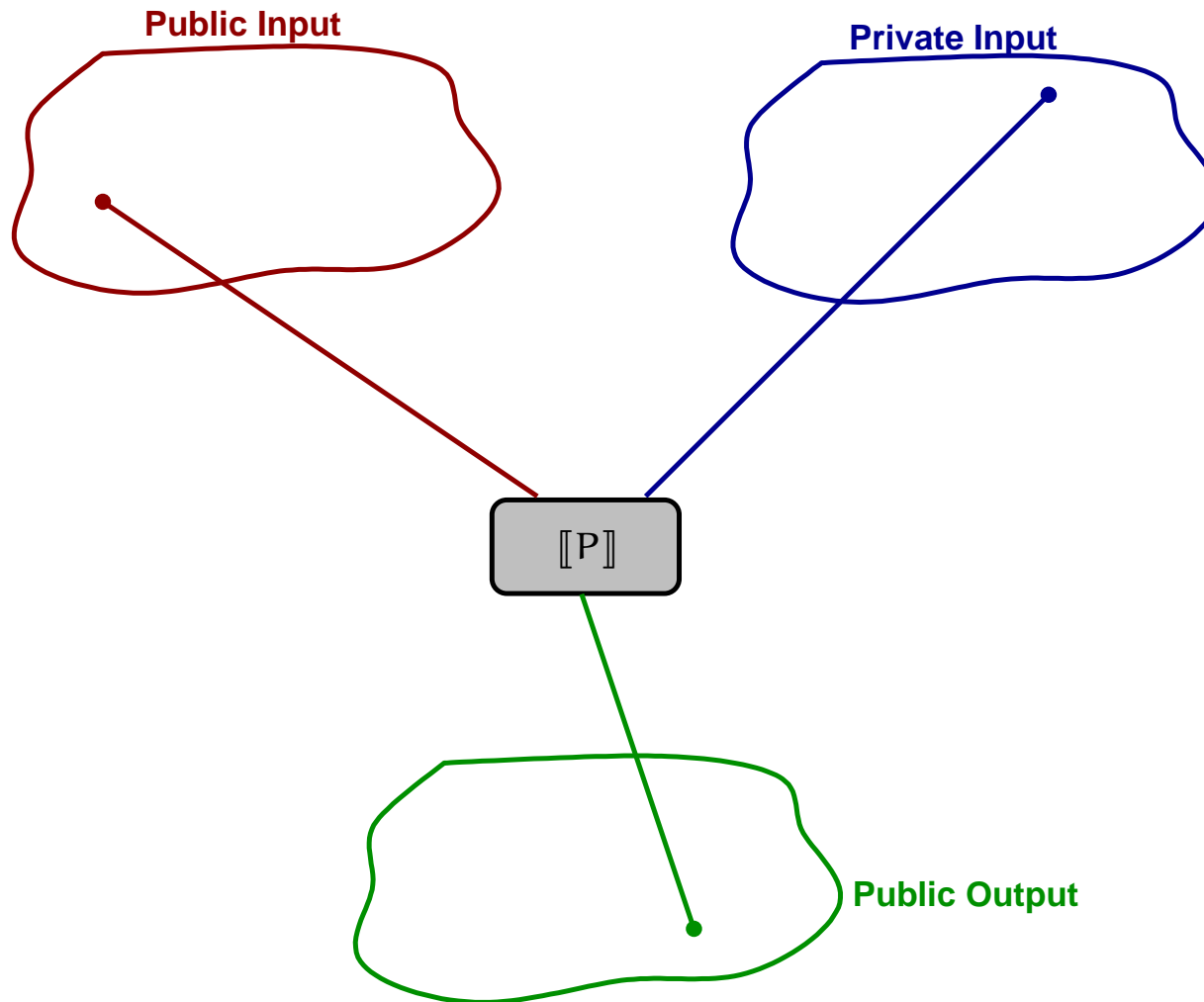
$$\forall l : L, \forall h_1, h_2 : H. [[P]](h_1, l)^L = [[P]](h_2, l)^L$$

Standard non-interference



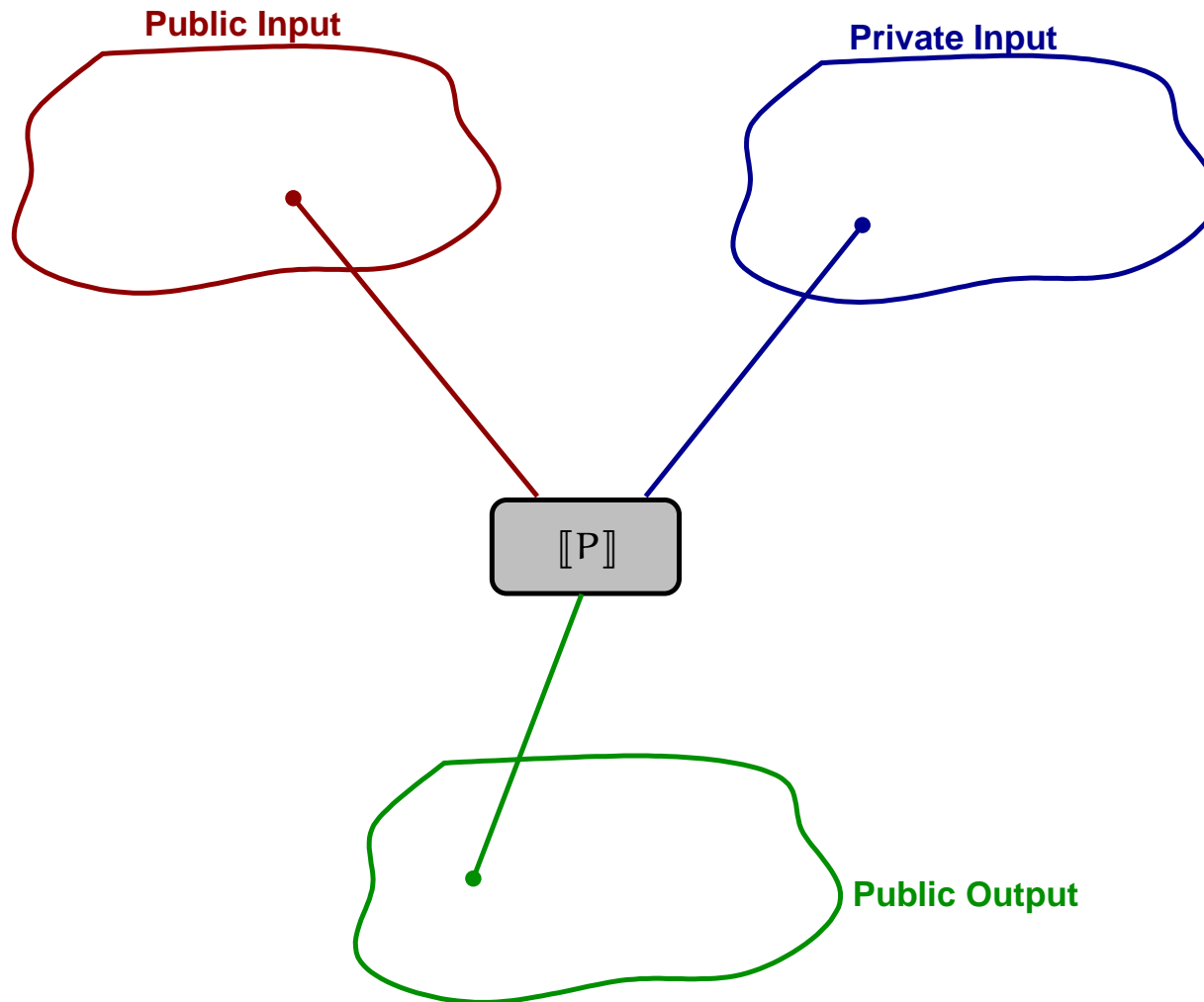
$$\forall l : L, \forall h_1, h_2 : H. \llbracket P \rrbracket(h_1, l)^L = \llbracket P \rrbracket(h_2, l)^L$$

Standard non-interference



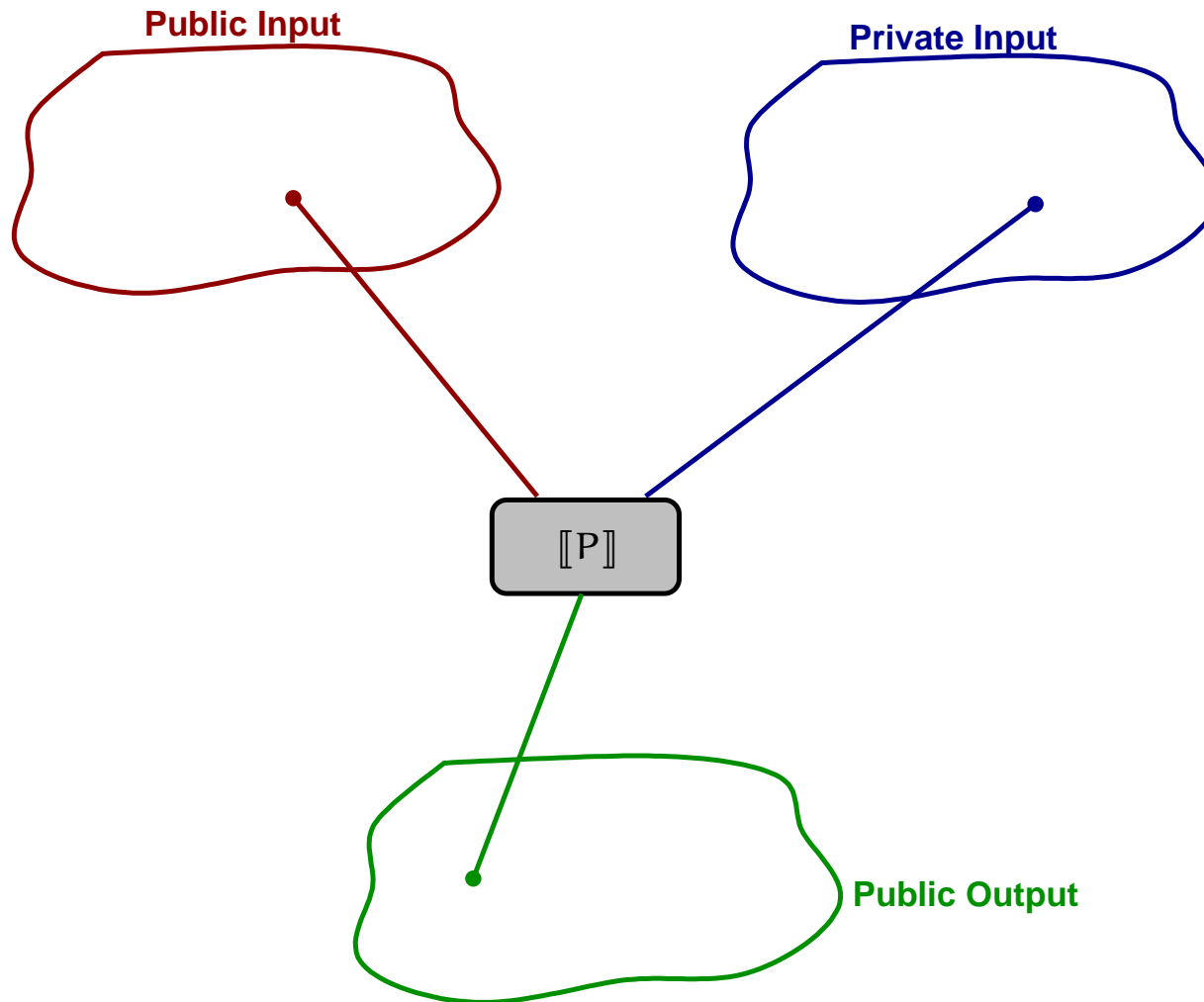
$$\forall l : L, \forall h_1, h_2 : H. [[P]](h_1, l)^L = [[P]](h_2, l)^L$$

Standard non-interference



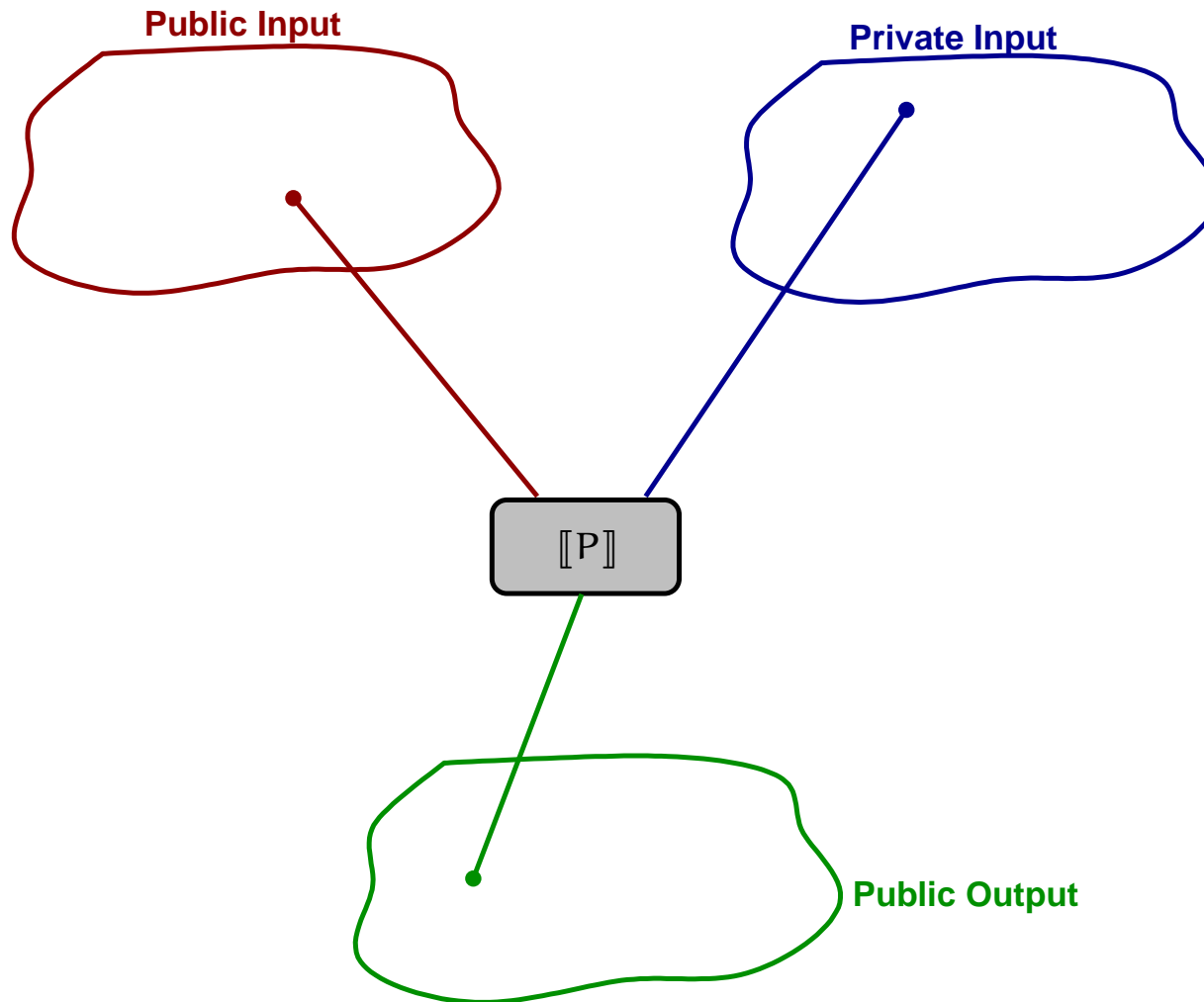
$$\forall l : L, \forall h_1, h_2 : H. [[P]](h_1, l)^L = [[P]](h_2, l)^L$$

Standard non-interference



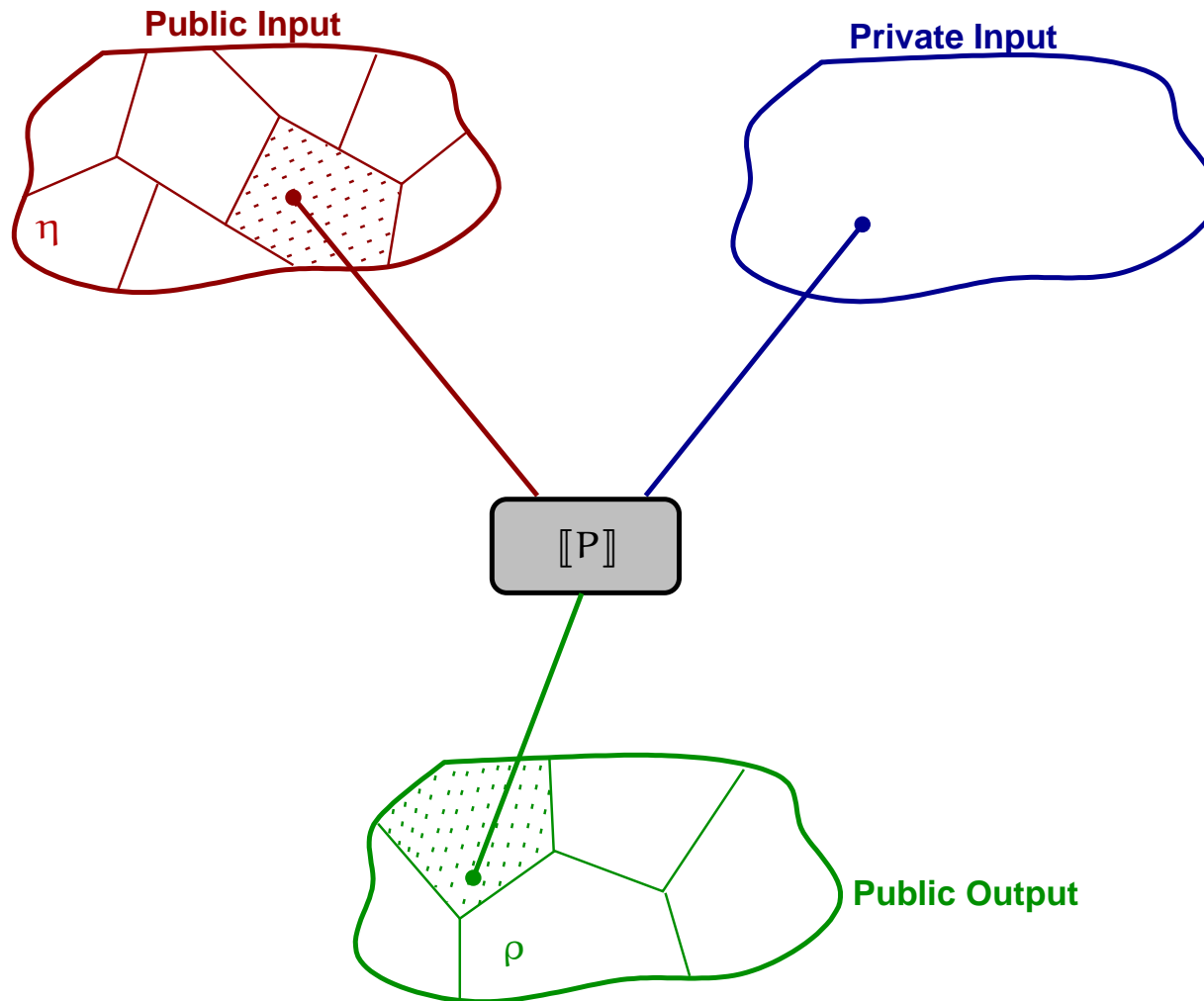
$$\forall l : L, \forall h_1, h_2 : H. [[P]](h_1, l)^L = [[P]](h_2, l)^L$$

Standard non-interference



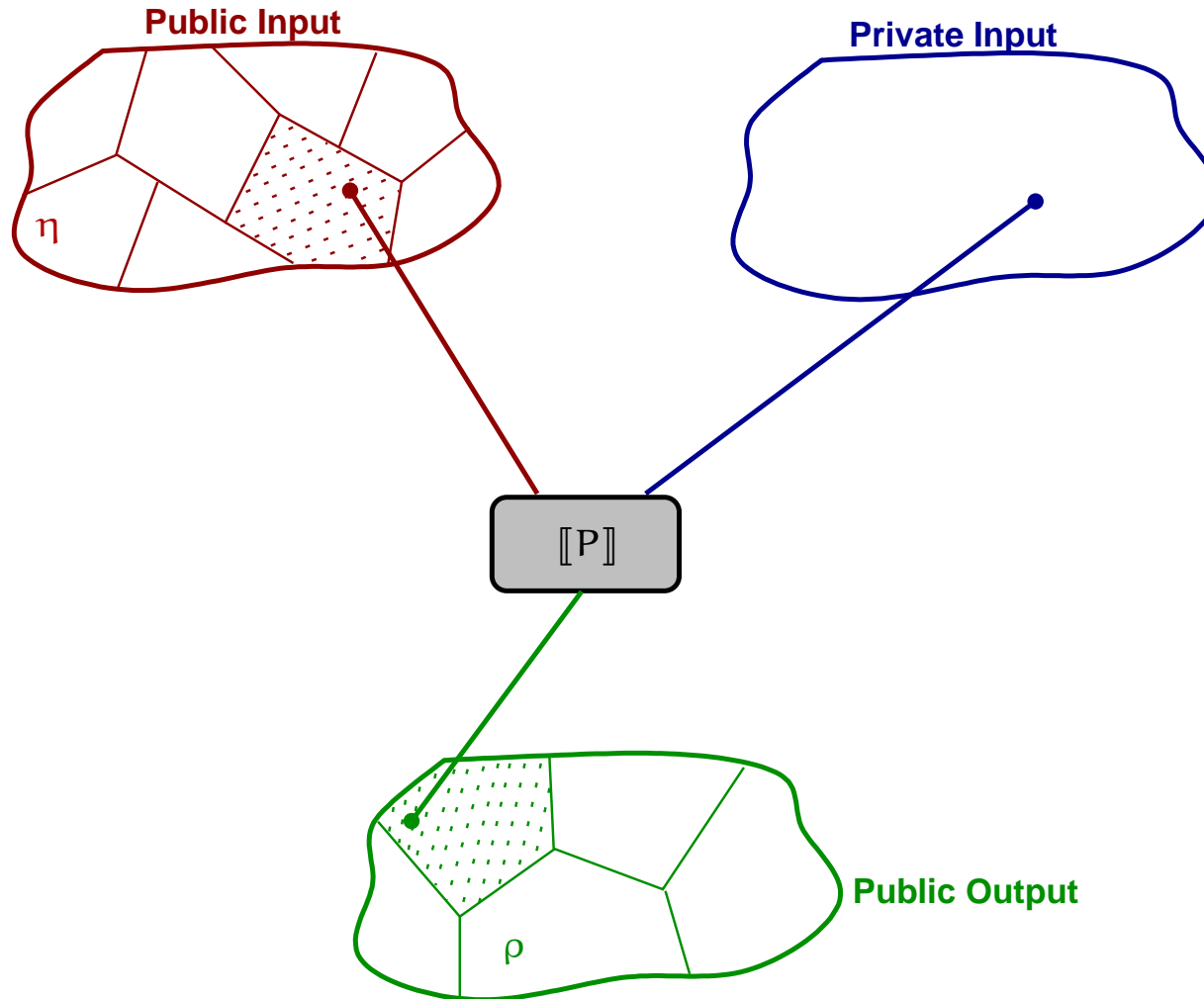
$$\forall l : L, \forall h_1, h_2 : H. [[P]](h_1, l)^L = [[P]](h_2, l)^L$$

Abstracting non-interference I: Narrow ANI



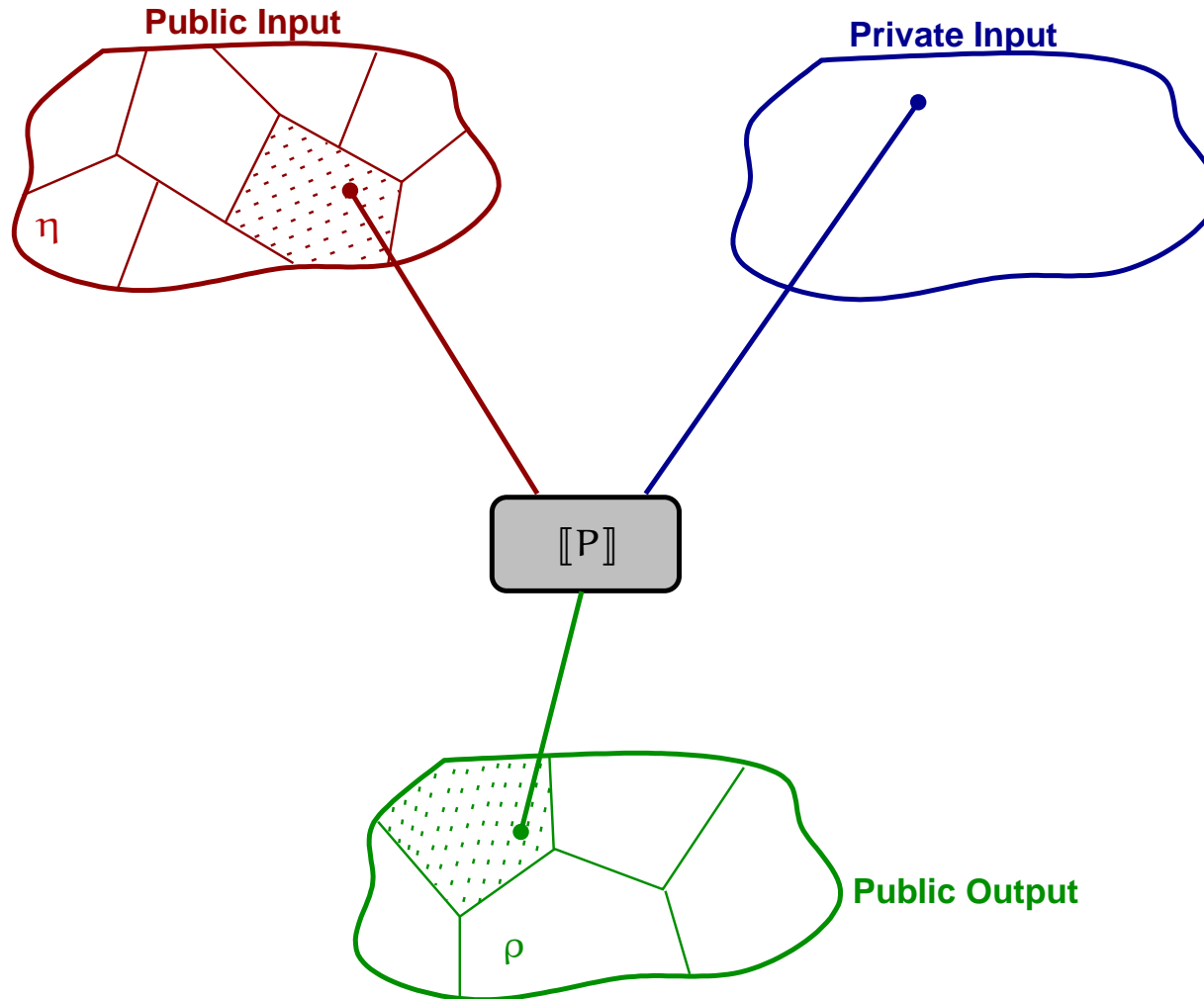
$$\rho, \eta \in \text{Abs}(\wp(\mathbb{V}^L)): [\eta]P(\rho): \eta(l_1) = \eta(l_2) \Rightarrow \rho([\![P]\!](h_1, l_1)^L) = \rho([\![P]\!](h_2, l_2)^L)$$

Abstracting non-interference I: Narrow ANI



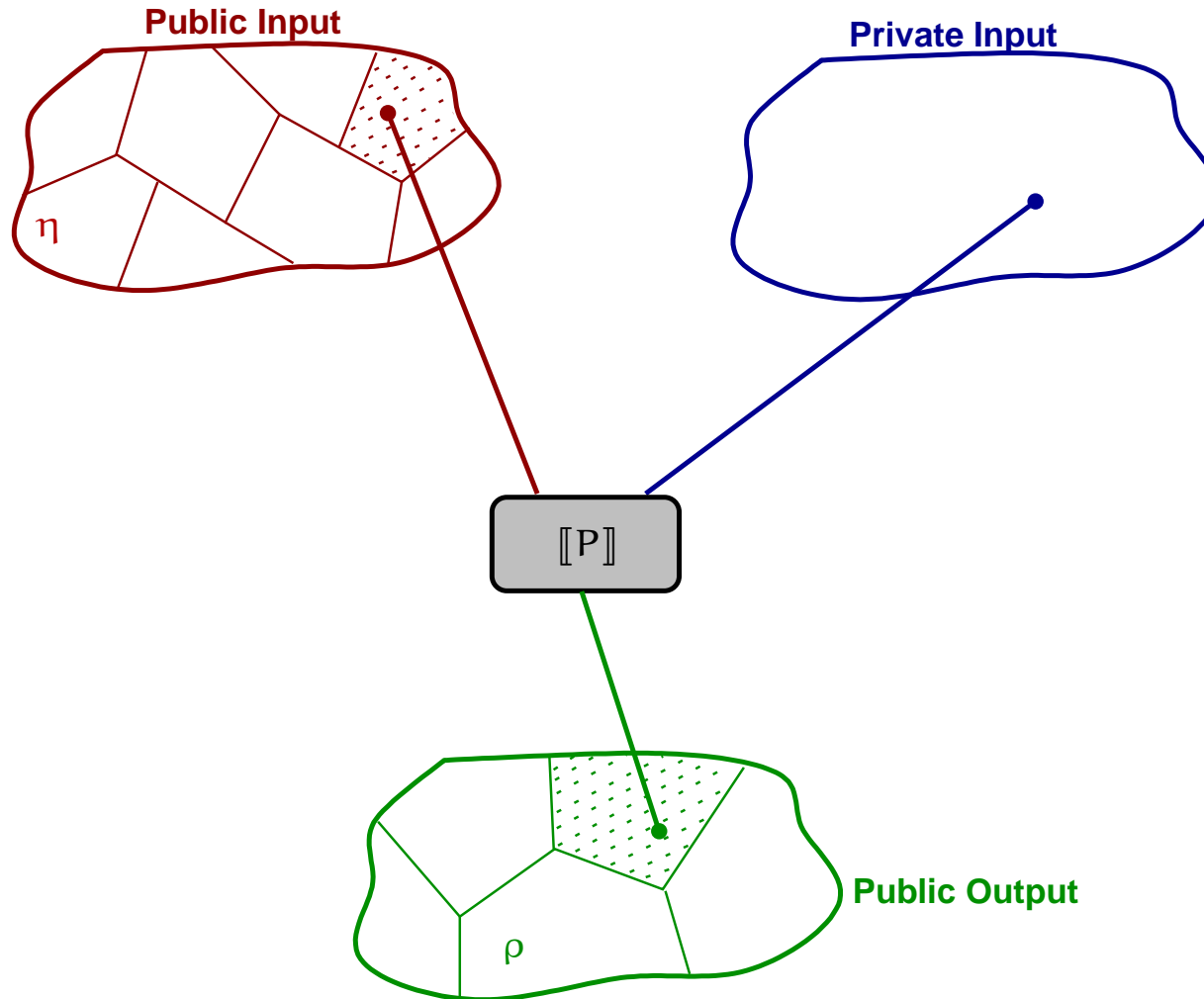
$$\rho, \eta \in \text{Abs}(\wp(\mathbb{V}^L)): [\eta]P(\rho): \eta(l_1) = \eta(l_2) \Rightarrow \rho([\![P]\!](h_1, l_1)^L) = \rho([\![P]\!](h_2, l_2)^L)$$

Abstracting non-interference I: Narrow ANI



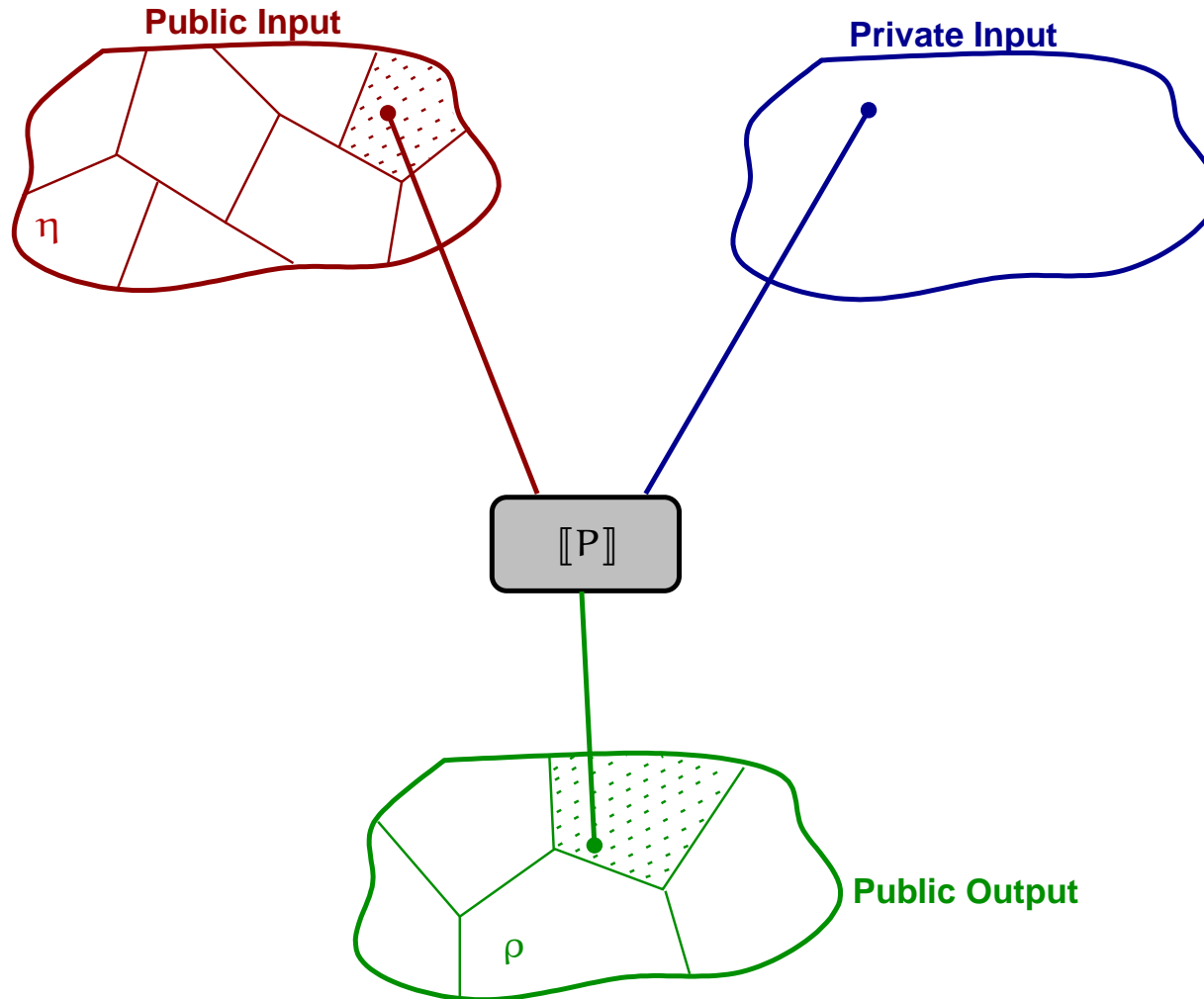
$$\rho, \eta \in \text{Abs}(\wp(\mathbb{V}^L)): [\eta]P(\rho): \eta(l_1) = \eta(l_2) \Rightarrow \rho([\![P]\!](h_1, l_1)^L) = \rho([\![P]\!](h_2, l_2)^L)$$

Abstracting non-interference I: Narrow ANI



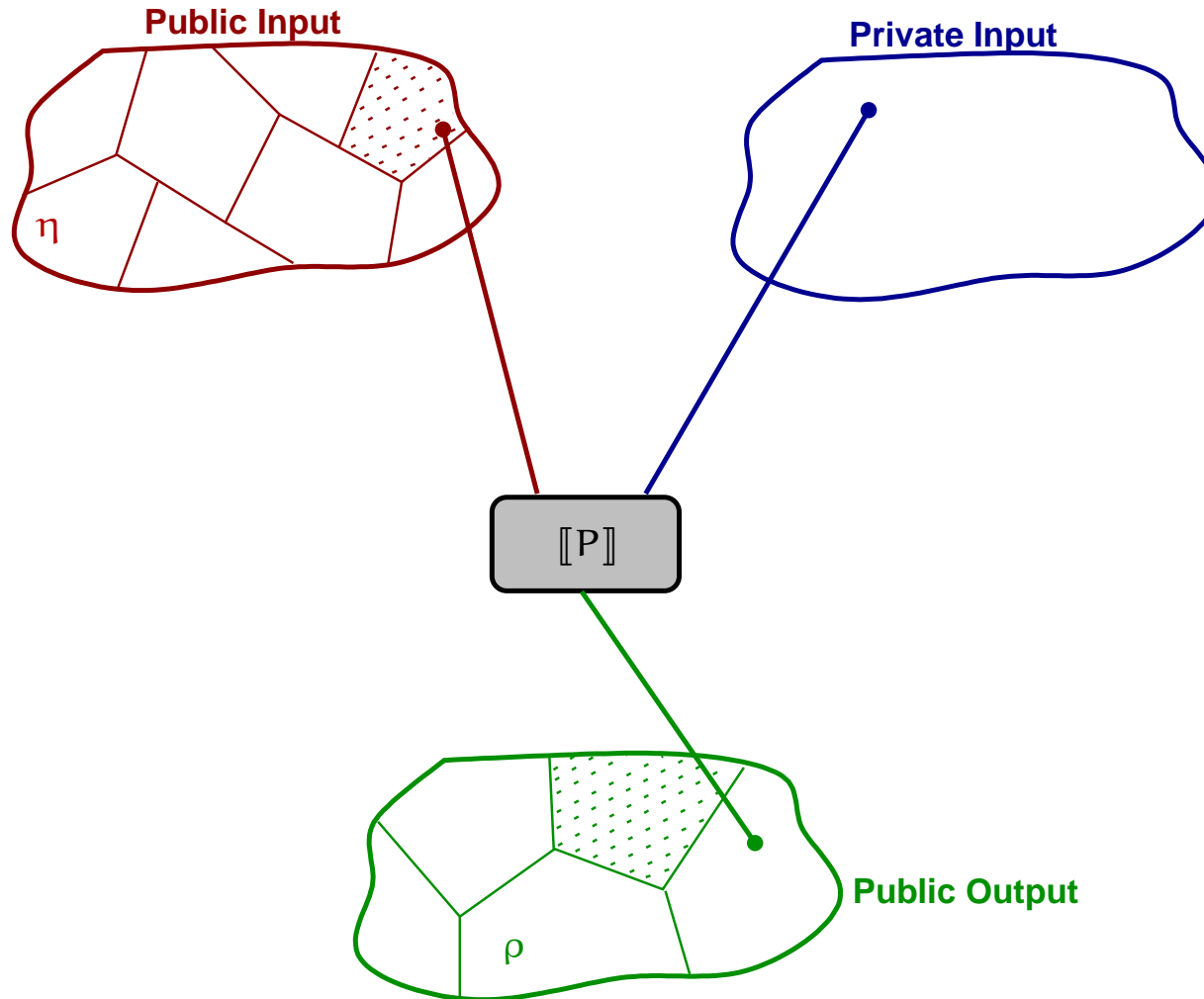
$$\rho, \eta \in \text{Abs}(\wp(\mathbb{V}^L)): [\eta]P(\rho): \eta(l_1) = \eta(l_2) \Rightarrow \rho([\![P]\!](h_1, l_1)^L) = \rho([\![P]\!](h_2, l_2)^L)$$

Abstracting non-interference I: Narrow ANI



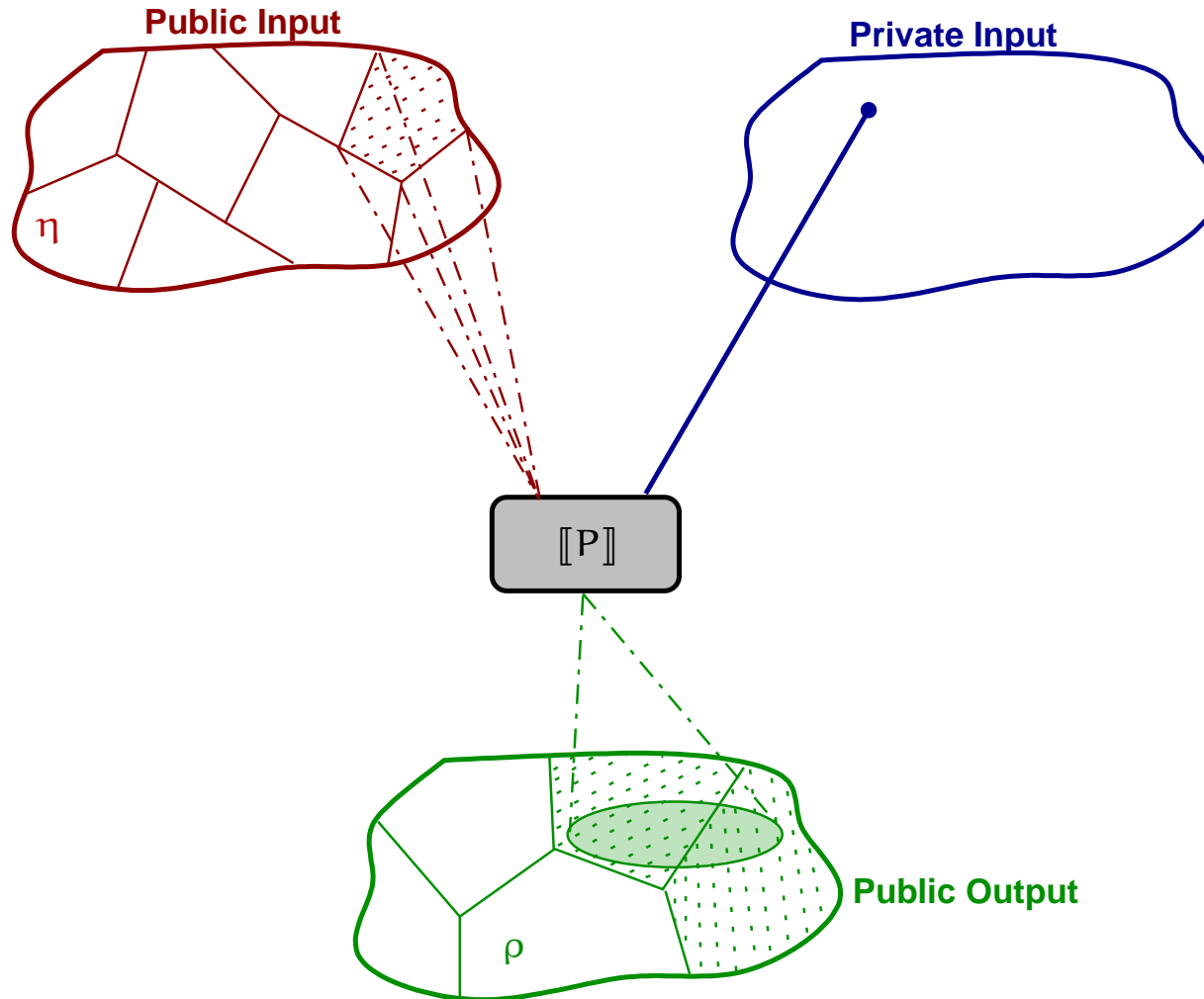
$$\rho, \eta \in \text{Abs}(\wp(\mathbb{V}^L)): [\eta]P(\rho): \eta(l_1) = \eta(l_2) \Rightarrow \rho([\![P]\!](h_1, l_1)^L) = \rho([\![P]\!](h_2, l_2)^L)$$

Abstracting non-interference I: Narrow ANI



$$\rho, \eta \in \text{Abs}(\wp(\mathbb{V}^L)): [\eta]P(\rho): \eta(l_1) = \eta(l_2) \Rightarrow \rho([\![P]\!](h_1, l_1)^L) = \rho([\![P]\!](h_2, l_2)^L)$$

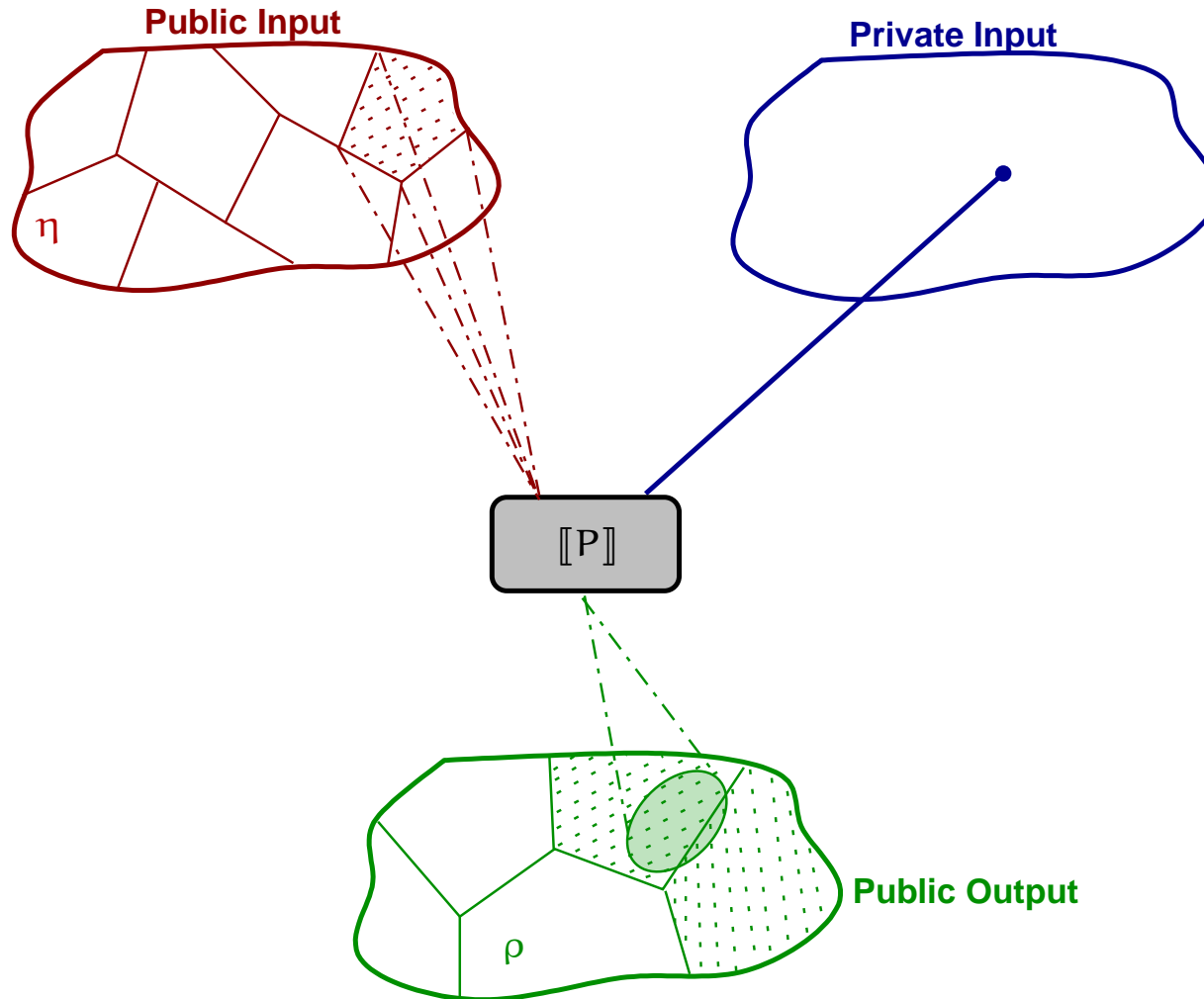
Abstracting non-interference II



$$\rho, \eta \in \text{Abs}(\wp(\mathbb{V}^L)) : (\eta)P(\rho) :$$

$$\eta(l_1) = \eta(l_2) \Rightarrow \rho(\llbracket P \rrbracket(h_1, \eta(l_1))^L) = \rho(\llbracket P \rrbracket(h_2, \eta(l_2))^L)$$

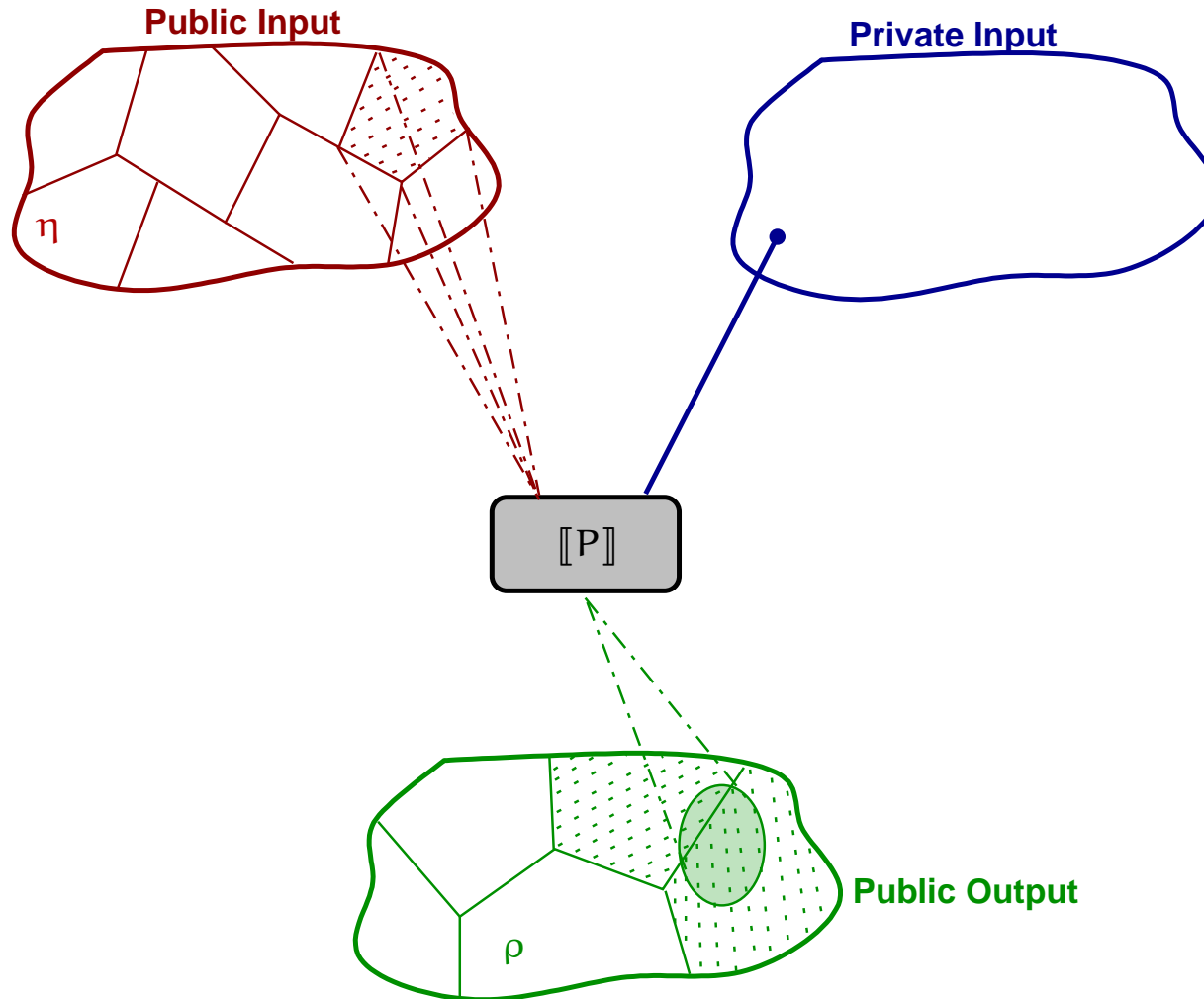
Abstracting non-interference II



$$\rho, \eta \in \text{Abs}(\wp(\mathbb{V}^L)) : (\eta)P(\rho) :$$

$$\eta(l_1) = \eta(l_2) \Rightarrow \rho(\llbracket P \rrbracket(h_1, \eta(l_1))^L) = \rho(\llbracket P \rrbracket(h_2, \eta(l_2))^L)$$

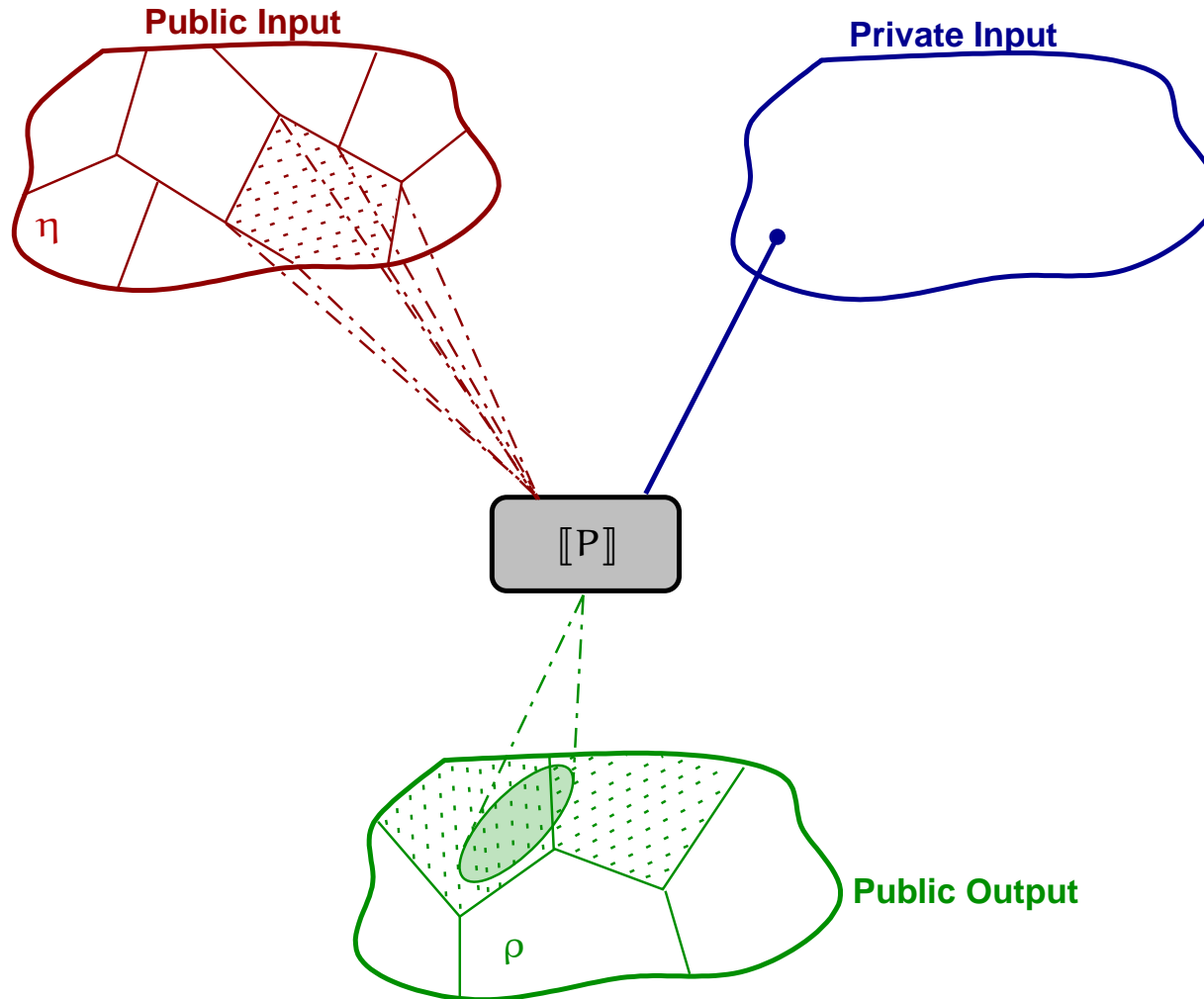
Abstracting non-interference II



$$\rho, \eta \in \text{Abs}(\wp(\mathbb{V}^L)) : (\eta)P(\rho) :$$

$$\eta(l_1) = \eta(l_2) \Rightarrow \rho(\llbracket P \rrbracket(h_1, \eta(l_1))^L) = \rho(\llbracket P \rrbracket(h_2, \eta(l_2))^L)$$

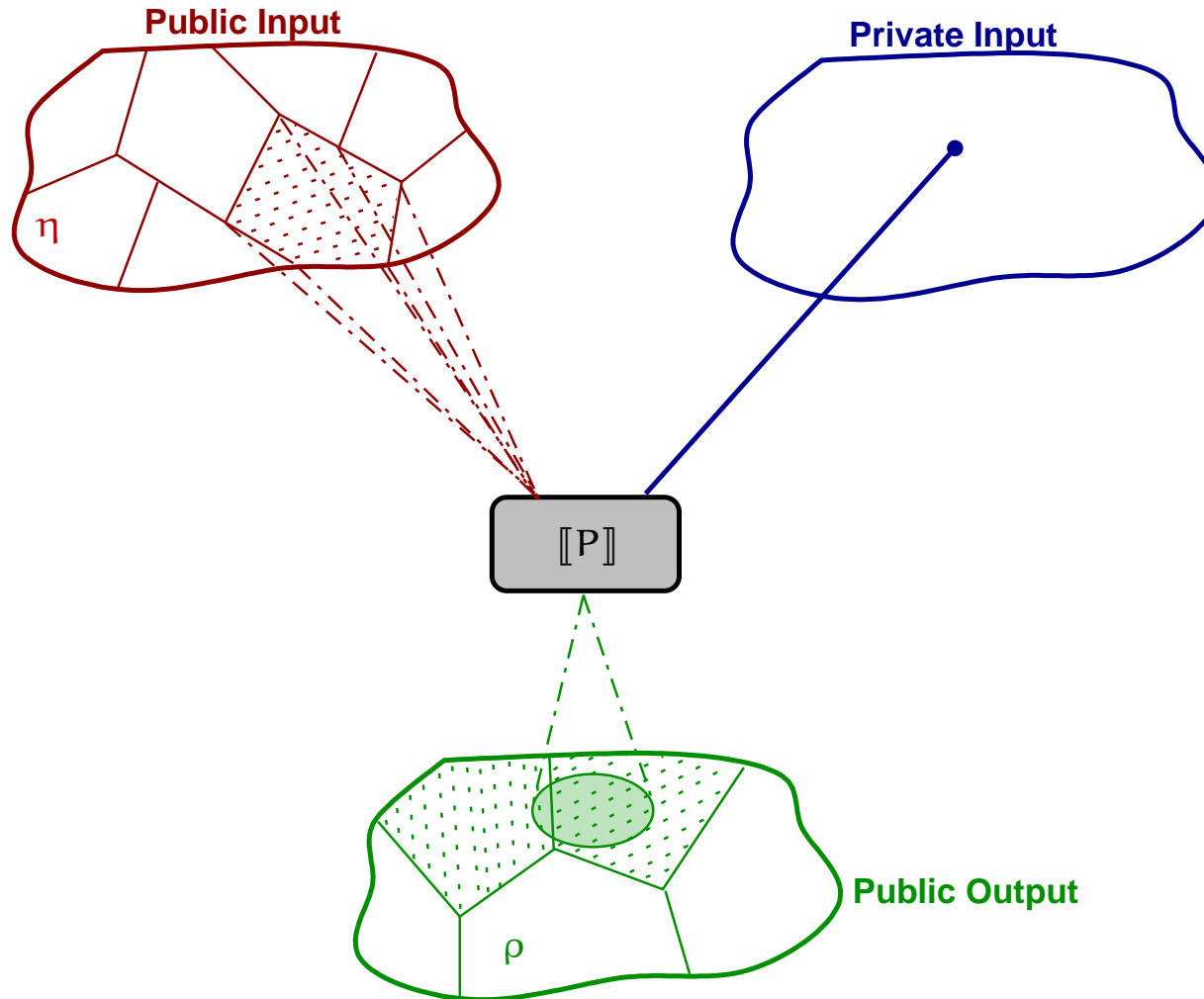
Abstracting non-interference II



$$\rho, \eta \in \text{Abs}(\wp(\mathbb{V}^L)) : (\eta)P(\rho) :$$

$$\eta(l_1) = \eta(l_2) \Rightarrow \rho(\llbracket P \rrbracket(h_1, \eta(l_1))^L) = \rho(\llbracket P \rrbracket(h_2, \eta(l_2))^L)$$

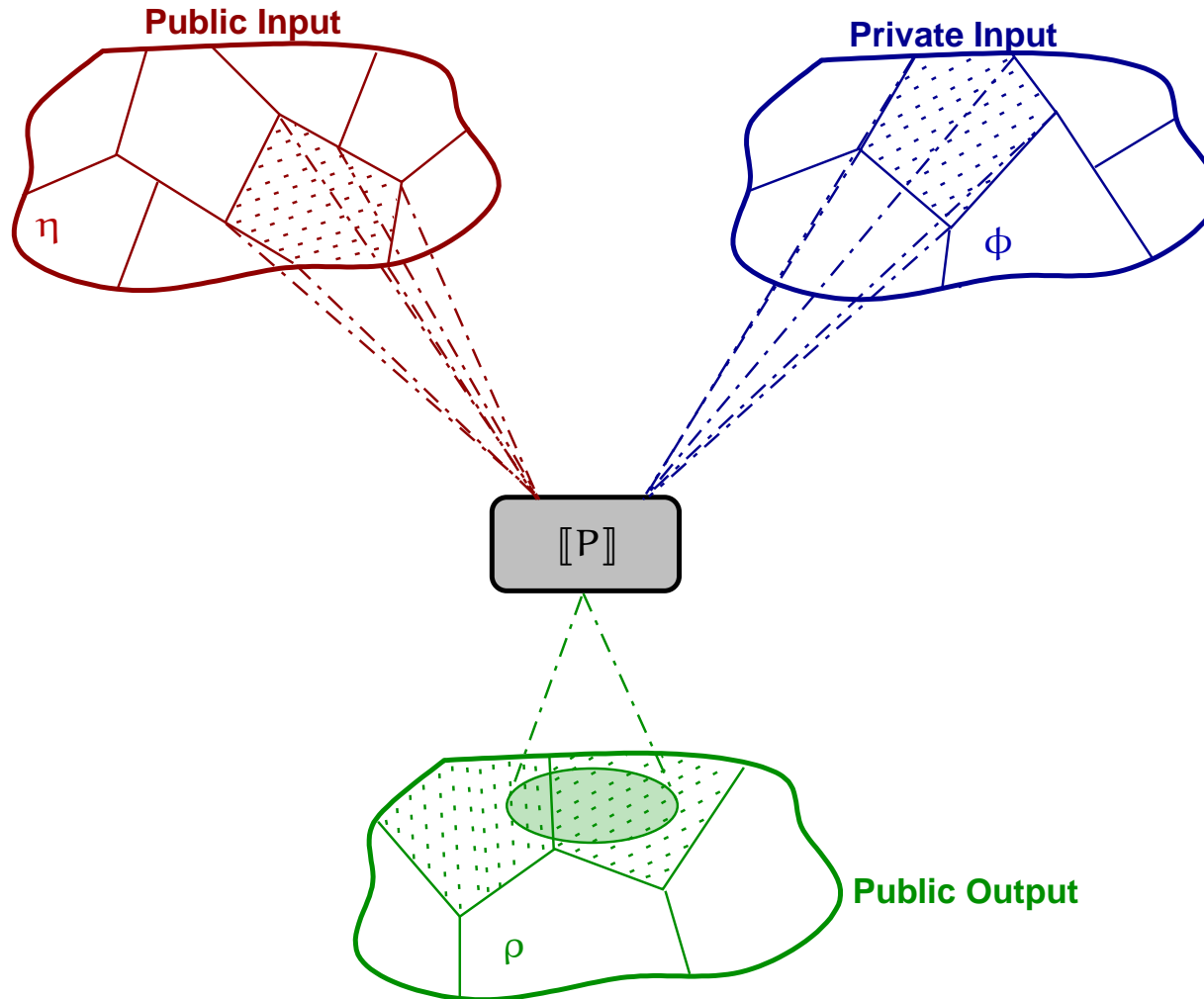
Abstracting non-interference II



$$\rho, \eta \in \text{Abs}(\wp(\mathbb{V}^L)) : (\eta)P(\rho) :$$

$$\eta(l_1) = \eta(l_2) \Rightarrow \rho(\llbracket P \rrbracket(h_1, \eta(l_1))^L) = \rho(\llbracket P \rrbracket(h_2, \eta(l_2))^L)$$

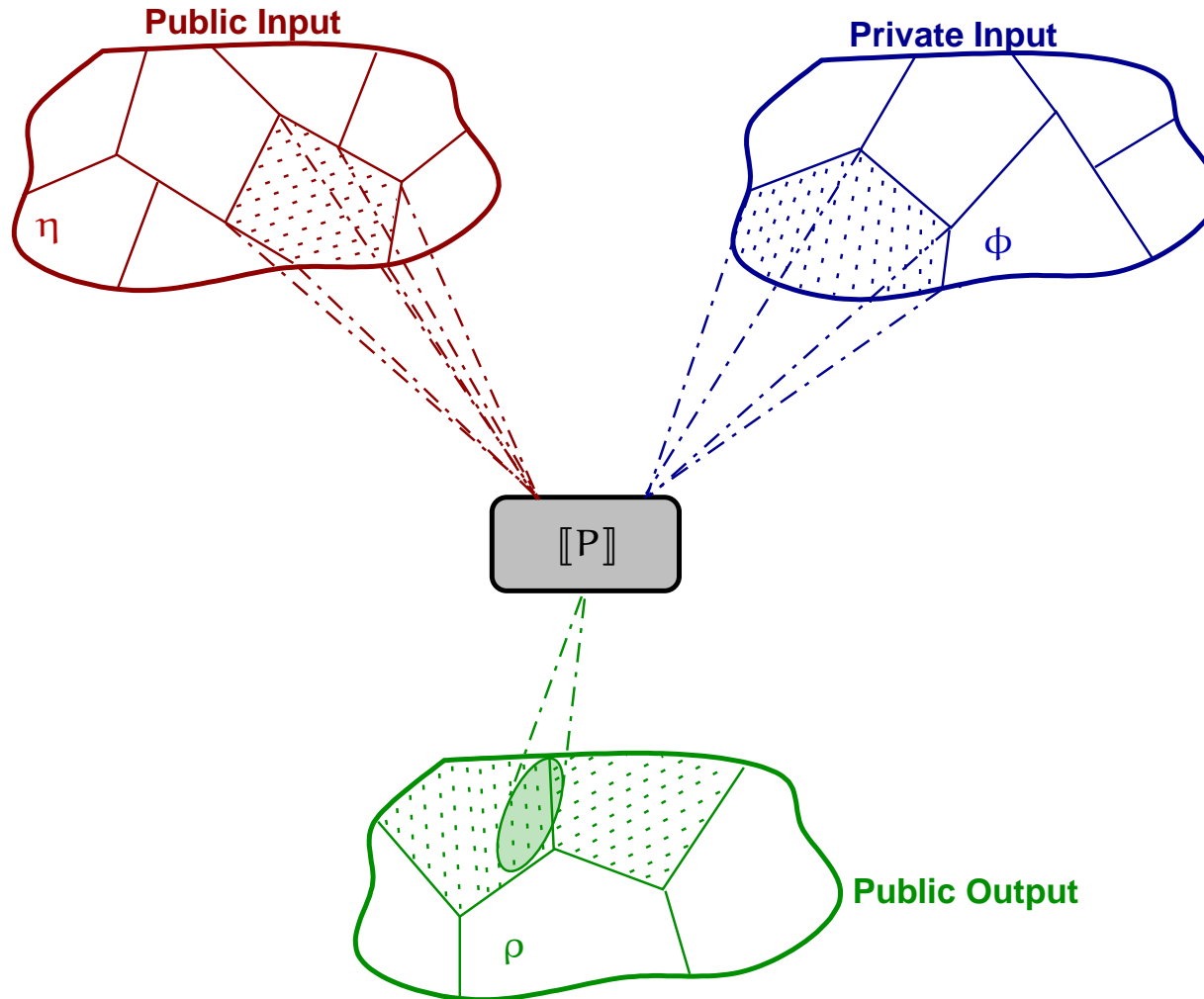
Abstracting non-interference III: ANI



$$\rho, \eta \in \text{Abs}(\wp(\mathbb{V}^L)), \phi \in \text{Abs}(\wp(\mathbb{V}^H)): (\eta)P(\phi \rightsquigarrow \rho):$$

$$\eta(l_1) = \eta(l_2) \Rightarrow \rho([P](\phi(h_1), \eta(l_1))^L) = \rho([P](\phi(h_2), \eta(l_2))^L)$$

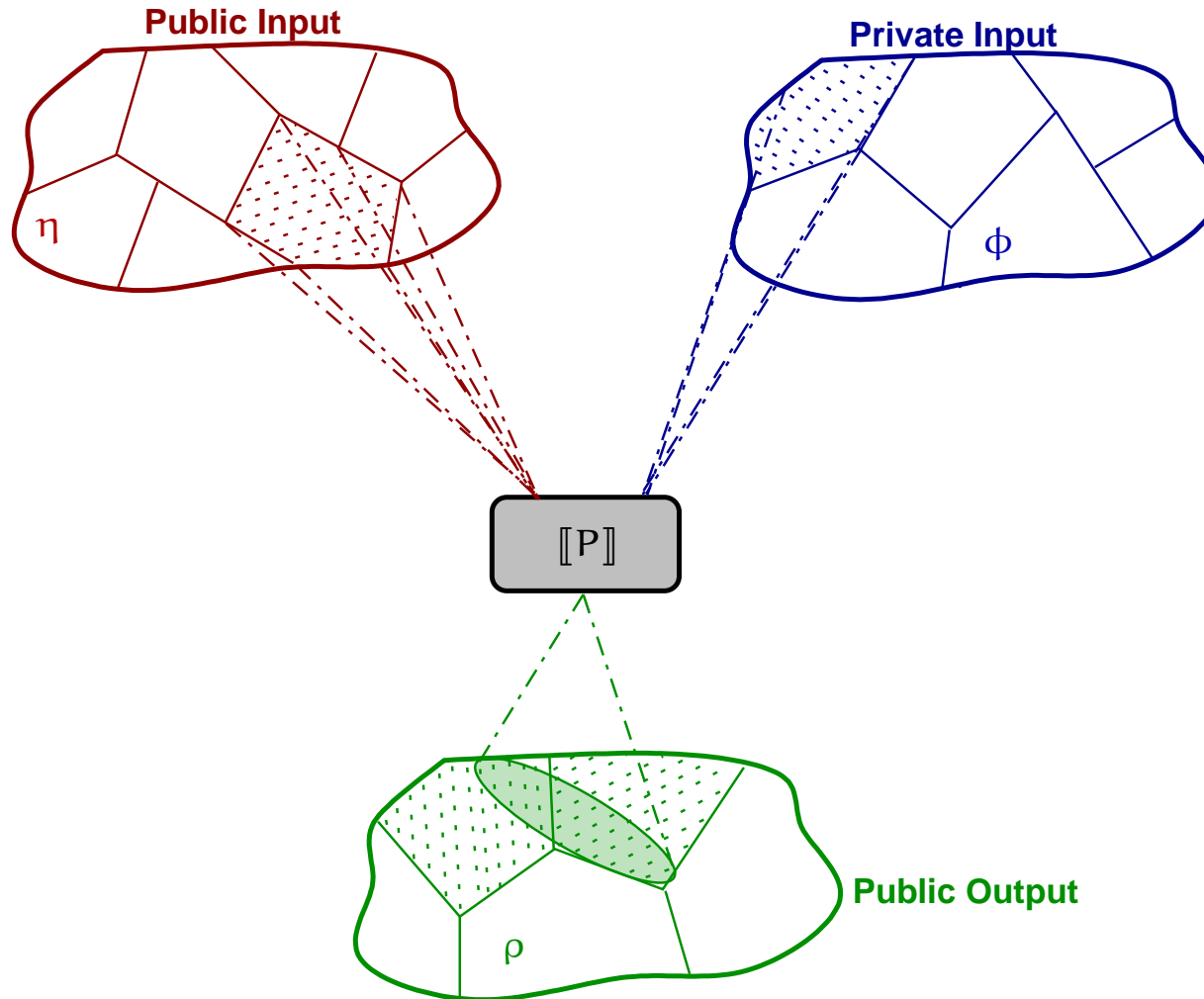
Abstracting non-interference II: ANI



$$\rho, \eta \in \text{Abs}(\wp(\mathbb{V}^L)), \phi \in \text{Abs}(\wp(\mathbb{V}^H)): (\eta)P(\phi \rightsquigarrow \rho):$$

$$\eta(l_1) = \eta(l_2) \Rightarrow \rho([P](\phi(h_1), \eta(l_1))^L) = \rho([P](\phi(h_2), \eta(l_2))^L)$$

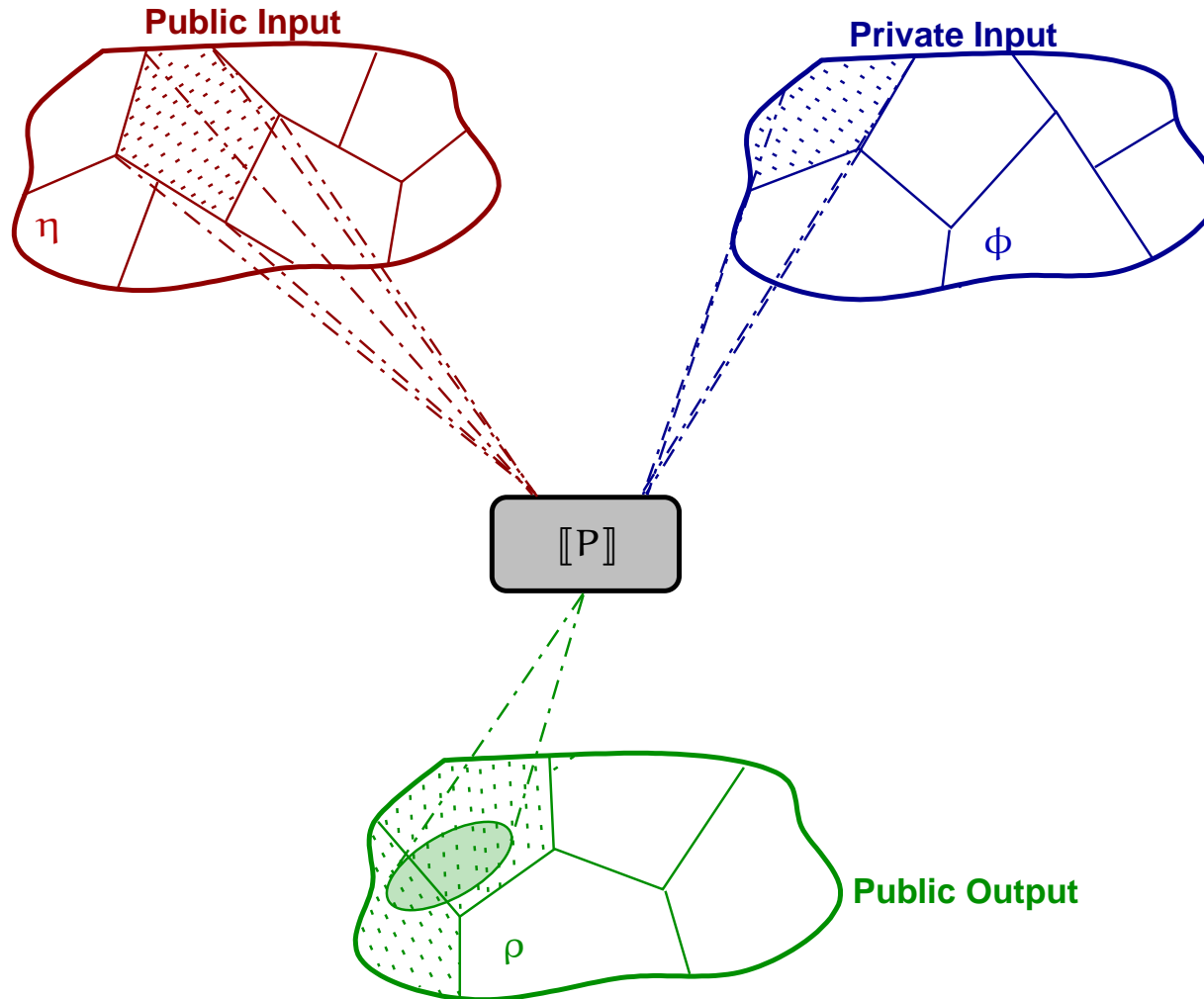
Abstracting non-interference III: ANI



$$\rho, \eta \in \text{Abs}(\wp(\mathbb{V}^L)), \phi \in \text{Abs}(\wp(\mathbb{V}^H)): (\eta)P(\phi \rightsquigarrow \rho):$$

$$\eta(l_1) = \eta(l_2) \Rightarrow \rho([P](\phi(h_1), \eta(l_1))^L) = \rho([P](\phi(h_2), \eta(l_2))^L)$$

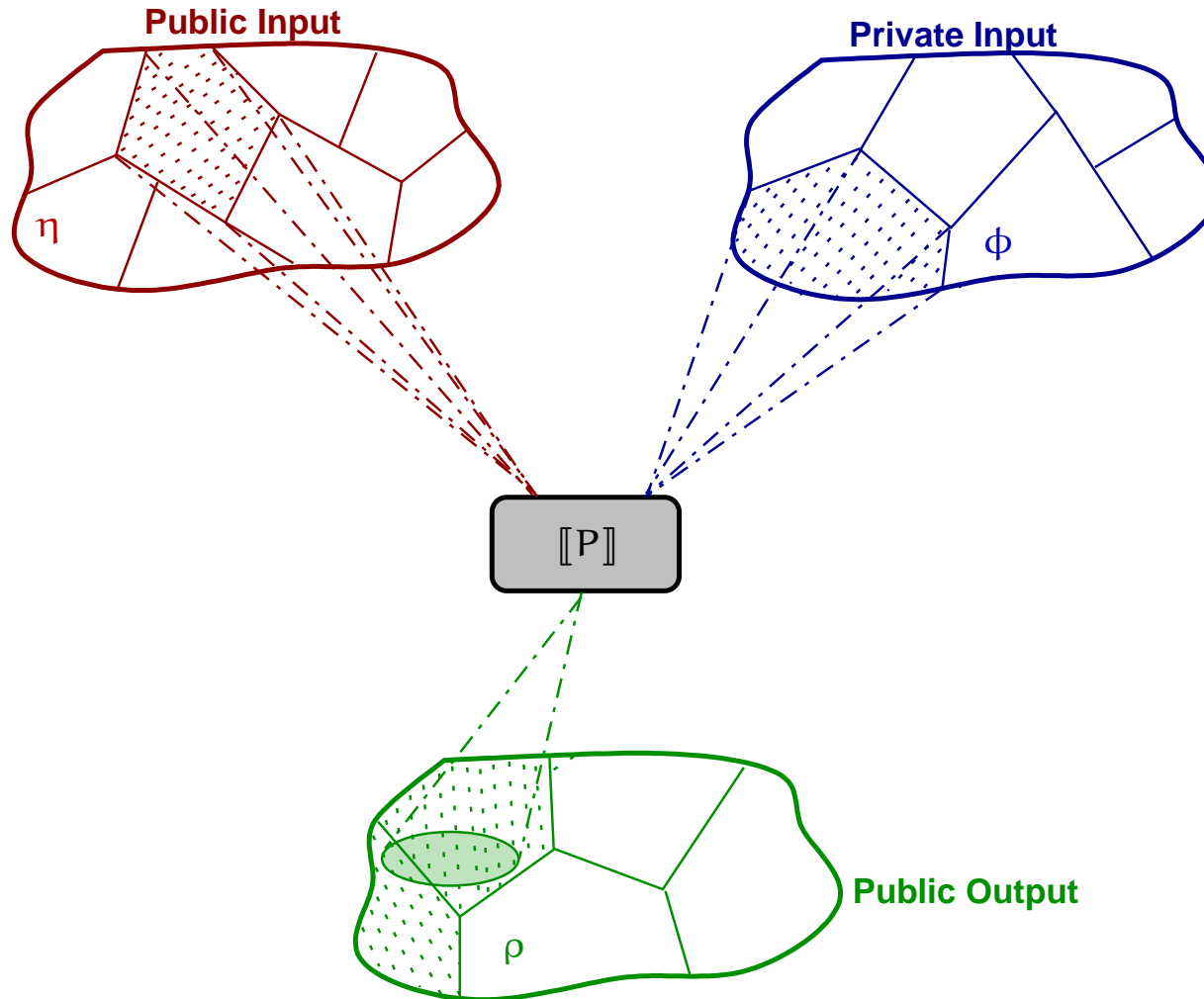
Abstracting non-interference III: ANI



$$\rho, \eta \in \text{Abs}(\wp(\mathbb{V}^L)), \phi \in \text{Abs}(\wp(\mathbb{V}^H)): (\eta)P(\phi \rightsquigarrow \rho):$$

$$\eta(l_1) = \eta(l_2) \Rightarrow \rho([P](\phi(h_1), \eta(l_1))^L) = \rho([P](\phi(h_2), \eta(l_2))^L)$$

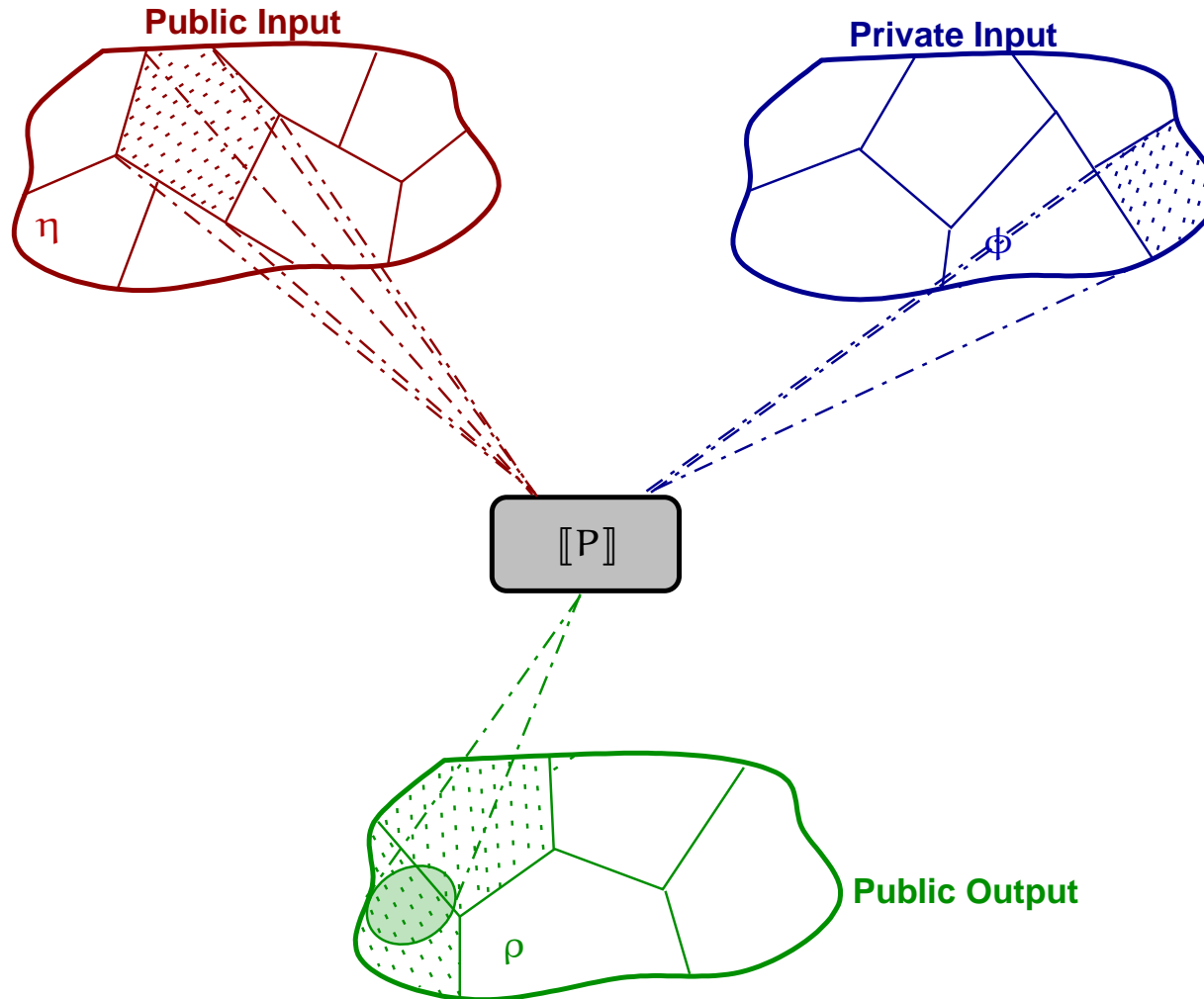
Abstracting non-interference III: ANI



$$\rho, \eta \in \text{Abs}(\wp(\mathbb{V}^L)), \phi \in \text{Abs}(\wp(\mathbb{V}^H)): (\eta)P(\phi \rightsquigarrow \rho):$$

$$\eta(l_1) = \eta(l_2) \Rightarrow \rho([P](\phi(h_1), \eta(l_1))^L) = \rho([P](\phi(h_2), \eta(l_2))^L)$$

Abstracting non-interference III: ANI



$$\rho, \eta \in \text{Abs}(\wp(\mathbb{V}^L)), \phi \in \text{Abs}(\wp(\mathbb{V}^H)): (\eta)P(\phi \rightsquigarrow \rho):$$

$$\eta(l_1) = \eta(l_2) \Rightarrow \rho([P](\phi(h_1), \eta(l_1))^L) = \rho([P](\phi(h_2), \eta(l_2))^L)$$

Deriving output attackers

Abstract interpretation provides advanced methods for designing abstractions
(refinement, simplification, compression ...) [Giacobazzi & Ranzato '97]

Designing abstractions = designing attackers

Deriving output attackers

Abstract interpretation provides advanced methods for designing abstractions (refinement, simplification, compression ...) [Giacobazzi & Ranzato '97]

Designing abstractions = designing attackers



- ⑥ Characterize the most concrete ρ such that $(\eta)P(\phi \rightsquigarrow \rho)$
[The most powerful *public observer*]

Deriving output attackers

Abstract interpretation provides advanced methods for designing abstractions (refinement, simplification, compression ...) [Giacobazzi & Ranzato '97]

Designing abstractions = designing attackers



- ⑥ Characterize the most concrete ρ such that $(\eta)P(\phi \rightsquigarrow \rho)$
[The most powerful *public observer*]

⇒ This would provide a certificate for security with a fixed input observation.

Timed abstract non-interference...

Standard denotational semantics

$\llbracket P \rrbracket$

$\alpha^{\mathcal{D}}$

Standard trace semantics

$\langle P \rangle$

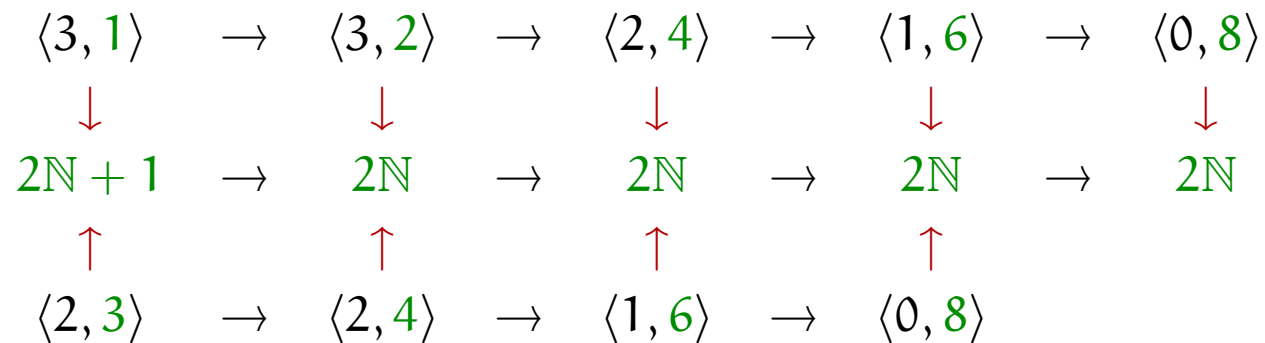
Trace semantics



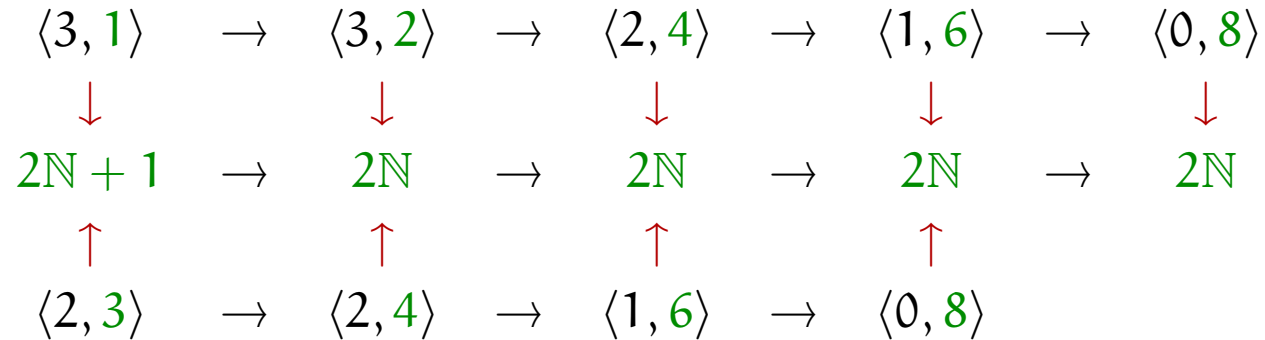
Traces' length = TIME ELAPSED

On traces

```
l := l + 1;  
while h ≠ 0 do  
    h := h - 1; l := l + 2;
```



On traces



Stuttering removes **time** from **traces**!

TRACE SEMANTICS

$$\begin{array}{c} 2\mathbb{Z} \longrightarrow 2\mathbb{Z} \longrightarrow 2\mathbb{Z} + 1 \longrightarrow 2\mathbb{Z} \\ \neq \\ 2\mathbb{Z} \longrightarrow 2\mathbb{Z} + 1 \longrightarrow 2\mathbb{Z} + 1 \longrightarrow 2\mathbb{Z} \end{array}$$

TRACE WITHOUT STUTTERING

$$\begin{array}{c} 2\mathbb{Z} \longrightarrow 2\mathbb{Z} + 1 \longrightarrow 2\mathbb{Z} \\ = \\ 2\mathbb{Z} \longrightarrow 2\mathbb{Z} + 1 \longrightarrow 2\mathbb{Z} \end{array}$$

Timed abstract non-interference...

Standard denotational semantics

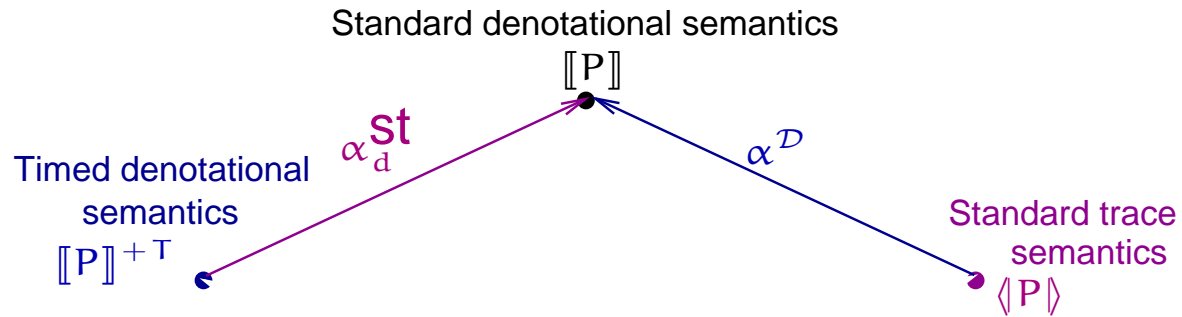
$\llbracket P \rrbracket$

α^D

Standard trace semantics

$\langle P \rangle$

Timed abstract non-interference...

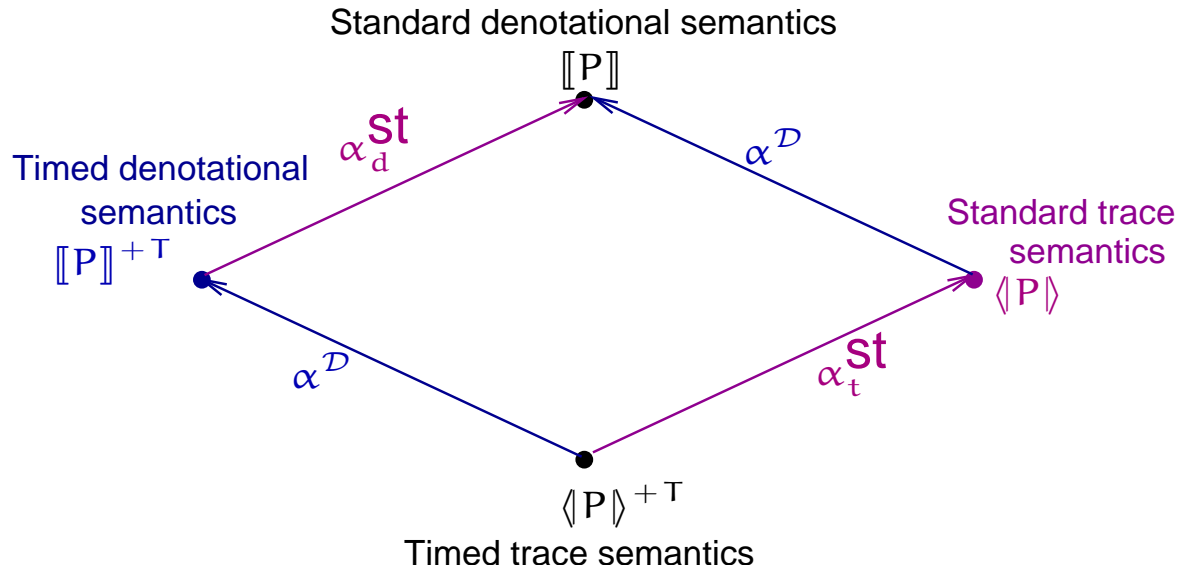


Timed denotational semantics



Time counter = TIME ELAPSED

Timed abstract non-interference...



Timed denotational semantics



Time counter = TIME ELAPSED

On timed semantics

$$t_A = t_T = 1$$

```
h := h mod 4;  
while h ≠ 0 do  
    h := h - 1; l := 2l - l;
```

$$\begin{array}{ccccccccc} \langle 7, 2, 0 \rangle & \rightarrow & \langle 3, 2, 1 \rangle & \rightarrow & \langle 2, 2, 4 \rangle & \rightarrow & \langle 1, 2, 7 \rangle & \rightarrow & \langle 0, 2, 10 \rangle \\ \langle 10, 2, 0 \rangle & \rightarrow & \langle 2, 2, 1 \rangle & \rightarrow & \langle 1, 2, 4 \rangle & \rightarrow & \langle 0, 2, 7 \rangle & & \end{array}$$

In general $\llbracket P \rrbracket(\langle h, l, 0 \rangle) = \langle 0, l, 3h + 1 \rangle$.

On timed semantics

$$t_A = t_T = 1$$

$$\begin{array}{ccccccccc} \langle 7, 2, 0 \rangle & \rightarrow & \langle 3, 2, 1 \rangle & \rightarrow & \langle 2, 2, 4 \rangle & \rightarrow & \langle 1, 2, 7 \rangle & \rightarrow & \langle 0, 2, 10 \rangle \\ \langle 10, 2, 0 \rangle & \rightarrow & \langle 2, 2, 1 \rangle & \rightarrow & \langle 1, 2, 4 \rangle & \rightarrow & \langle 0, 2, 7 \rangle & & \end{array}$$

In general $\llbracket P \rrbracket(\langle h, l, 0 \rangle) = \langle 0, l, 3h + 1 \rangle$.

\Rightarrow We can derive the maximal harmless observable property of time.

On timed semantics

$$t_A = t_T = 1$$

$$\begin{array}{ccccccc} \langle 7, 2, 0 \rangle & \rightarrow & \langle 3, 2, 1 \rangle & \rightarrow & \langle 2, 2, 4 \rangle & \rightarrow & \langle 1, 2, 7 \rangle & \rightarrow & \langle 0, 2, 10 \rangle \\ \langle 10, 2, 0 \rangle & \rightarrow & \langle 2, 2, 1 \rangle & \rightarrow & \langle 1, 2, 4 \rangle & \rightarrow & \langle 0, 2, 7 \rangle & & \end{array}$$

In general $\llbracket P \rrbracket(\langle h, l, 0 \rangle) = \langle 0, l, 3h + 1 \rangle$.

We provide necessary and sufficient conditions such that



Abstraction removes **time** from **timed denotational semantics**!

TIMED SEMANTICS

$$\text{Par}(\llbracket P \rrbracket(2, 4, 0))^{TL} = \text{Par}(6, 3) = \langle 2\mathbb{Z}, 2\mathbb{Z} + 1 \rangle$$

\neq

$$\text{Par}(\llbracket P \rrbracket(4, 4, 0))^{TL} = \text{Par}(8, 6) = \langle 2\mathbb{Z}, 2\mathbb{Z} \rangle$$

UNTIMED SEMANTICS

$$\Pi^T(\langle 2\mathbb{Z}, 2\mathbb{Z} + 1 \rangle) = \langle 2\mathbb{Z}, \mathbb{Z} \rangle$$

$=$

$$\Pi^T(\langle 2\mathbb{Z}, 2\mathbb{Z} \rangle) = \langle 2\mathbb{Z}, \mathbb{Z} \rangle$$

Conclusion

- ⑥ We extended the notion of abstract non-interference with time:
 - ▣ We can model stronger attackers able to measure time;
 - ▣ We can model abstract non-interference in systems with time;

Conclusion

- ⑥ We extended the notion of abstract non-interference with time:
 - ▣ We can model stronger attackers able to measure time;
 - ▣ We can model abstract non-interference in systems with time;
- ⑥ We can simply generalize this construction for systems where time is changed by time;

Conclusion

- ⑥ We extended the notion of abstract non-interference with time:
 - ▣ We can model stronger attackers able to measure time;
 - ▣ We can model abstract non-interference in systems with time;
- ⑥ We can simply generalize this construction for systems where time is changed by time;
- ⑥ We are working on a implementation of a tool for deriving certifications of abstract non-interference properties.