

25 Years of Abstract Interpretation

The German Perspective

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Who?

Reinhard **Wilhelm** Saarbrücken

Bernhard **Steffen** Dortmund

Jens **Knoop** Wien

Markus **Müller-Olm** Münster

Helmut **Seidl** München

Andreas **Podelski** Freiburg

What?

1. Grammar flow analysis
2. Worst-Case Execution Time (WCET)
3. Data flow analysis as model checking
4. Partial Redundancy Elimination (PRE)
5. Decidability frontiers for abstraction
6. Set-based analysis, set constraints
7. Abstraction refinement-based model checking (ARMC)
8. Transition invariants, termination analysis

I. Grammar Flow Analysis

- abstract one-step derivation (“ $N \rightarrow E$ ”)
- **concrete** lattice of sets of derivation trees,
abstract lattice of graphs
- first/follow, emptiness,
“restraints” between attributes, ...
- **1982**: Möncke & Wilhelm (2006: C & C)

2. Worst-Case Execution Time (WCET)

- sum up time that all instructions can take on path
(take maximum over all execution paths)
- cost of instruction depends on state

predicted cache hit \Rightarrow tighter WCET bound

- analysis of reachable configurations: **cache**, pipelines, ...
- SAS'96: Wilhelm, Ferdinand, Martin, ...
company: AbsInt (Airbus, Bosch, ...)

3. Data Flow Analysis as Model Checking

- data flow facts
 - = **temporal-logic** properties
 - = fixpoint expressions in mu-calculus
- fixpoints in finite abstract lattice? - model checking!
- conceptual, practical (company: MetaFrame)
- 1991: Steffen (1998: Schmidt, 2000: C & C)

4. Partial Redundancy Elimination (PRE)

- compiler optimization (code motion)
- extension of total redundancy (availability for all/some paths)
- rigorous proof of **correctness**
- **unidirectional**
- Knoop, Steffen, Rüthing, 1992 - ...

5. Decidability Frontiers for Abstraction

- When, or how far, do we **have to** abstract?
 - Ignore **guards!**
- What abstractions lead to decidability?
 - **Intervals!**
- 2002 - ... : Müller-Olm, Seidl

6. Set-based Analysis of Programs over Trees

- **set-based analysis** of program P over lists or other trees:
 1. transform $P \Rightarrow P^\#$ such that: $\text{post}_{P^\#} = \text{post of } P^\#$
 2. solve $P^\#$ ($P^\#$ is a **set constraint**)
- **set constraint = fixpoint equation** over tree grammars !
- *greatest solution characterizes non-termination*
- Reynolds'69, Jones'79, Heintze'89, C & C'92
- 1990 - ...: Frühwirth, Ganzinger, Seidl, Podelski, ...

7. Abstraction Refinement-based Model Checking

- **completeness** relative to widening
- ARMC: logic-based implementation
(*Andrey Rybalchenko's Model Checker*)
- abstraction by linear-arithmetic constraint solver (and theory extensions)
- verification of networked train control systems
- 2000 - ... : Veith, Podelski, Rybalchenko, ...

8. Transition Invariants, Termination Analysis

- abstract least fixpoint **not good** for termination !
- **transition invariant**
 - = summary
 - = abstract least fixpoint
- **transition invariant** can prove termination
 - ⇒ abstract least fixpoint **good** for termination !
- 2004 - ... : Cook, Podelski, Rybalchenko

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Ongoing

- Oldenburg - Saarbrücken - Freiburg
AVACS (DFG)
networked embedded systems
- Saarbrücken - Freiburg - Aachen MSRC - Redmond MSR
Verisoft (BMBF)
Microsoft Hypervisor
- München
PUMA (Graduiertenschule)
AI & types & theorem proving
- Saarbrücken
hierarchical shape analysis
- Münster, München
decidable abstractions of parallel programs
- Freiburg
symbolic shape analysis,
thread-modular analysis, ...
- ...

Challenges in AI

... in: Concurrency, Search, Control Theory

Concurrency

- state explosion
- thread-modular verification
... *Cartesian abstraction*
- synchronization
... *intersection of automata*
- confluence

AI for AI Planning

- abstraction for good estimates of distance between states
- abstraction is good if search is fast
- abstract interpretation for artificial intelligence?

Control Theory

- system with non-linear dynamics
(e.g. pendulum)
- safety
- stability

Towards a Conclusion ...

What is specific about German AI research?

Sets vs. Set Theory

- mathematicians **use sets**
(notation for formulating ideas)
- logicians **do set theory**
(study of the infinite)

AI vs. AI Theory

- program analysis researchers **use AI**
(notation for formulating ideas)
- semantics researchers **do AI theory**
(study of the infinite)
- German AI researchers are perhaps rather
among the **AI users**