

Homework.

Mathematical Logic - Gianluigi Bellin

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1(i). **Write derivations in the calculus of sequents** of the following expressions, by applying the procedure *semantic tableaux*

1. $A \rightarrow (B \rightarrow C), B \Rightarrow A \rightarrow C$;
2. $A \rightarrow B, B \rightarrow C \Rightarrow A \rightarrow C$;
3. $(A \rightarrow B) \rightarrow A \Rightarrow A$.

Consider a **Hilbert system** with axioms

Ax 1 $A \rightarrow (B \rightarrow A)$;

Ax 2 $(A \rightarrow (B \rightarrow C)) \rightarrow ((A \rightarrow B) \rightarrow (A \rightarrow C))$;

Ax 3 $\neg\neg A \rightarrow A$

and with Modus Ponens as only rule of inference:

$$\frac{A \rightarrow B \quad A}{B} \text{MP} .$$

Find derivations of

H1 $A \rightarrow (B \rightarrow C), B \vdash A \rightarrow C$;

H2 $A \rightarrow B, B \rightarrow C \vdash A \rightarrow C$

H3 $(A \rightarrow B) \rightarrow A \vdash A$

by applying the procedure in the proof of equivalence between Gentzen's systems and Hilbert systems.

Hints: In (H1) build a derivation $A \rightarrow (B \rightarrow C), B, A \vdash C$ and then apply the Deduction Theorem. You do not need to use Axiom 3. In (H3) you do need to use Axiom 3.

Remark. A *Natural Deduction* calculus is proof-system without axioms, with *introduction* and *elimination* inference rules for each connective and with rules that specify which sets of assumptions are *open* and which are *discharged* at each stage of the derivation. Proofs have the form of *trees with annotations* concerning the discharge of assumptions. Here *proof trees* are either an *open assumption* A or result from derivations d, d_1 and d_2 by the inferences

$$\begin{array}{c}
(1) \\
[A] \\
d \\
(1) \frac{B}{A \rightarrow B} \rightarrow\text{-intro} \\
[A] \text{ is a set of leaves labelled with } A. \\
\frac{\frac{d_1}{A \rightarrow B} \quad \frac{d_2}{A}}{B} \rightarrow \text{elim} \qquad \frac{d}{\neg\neg A} \neg\neg \text{ rule}
\end{array}$$

with the following *deduction rules*:

- (1) if $\Gamma, A \vdash B$ then $\Gamma \vdash A \rightarrow B$;
- (2) if $\Gamma \vdash A \rightarrow B$ and $\Delta \vdash A$ then $\Gamma, \Delta \vdash B$.
- (3) if $\Gamma \vdash \neg\neg A$ then $\Gamma \vdash A$.

Then one can prove in such a system the axioms for implication and negation in a Hilbert-style system and also give more perspicuous derivations of (H1)-(H3).