Logica Computazionale 2009 Exercise 2

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Exercise 1. Show that it is possible to falsify the following statement:

"If everybody loves someone, then someone is loved by everybody"

Hint: Construct a model $\mathcal{M} = (M, L)$, i.e., a set of individuals M and a two places relation L, such that L(x, y) is an interpretation of "x loves y" and also satisfies the first condition ("everybody likes someone"), but for a given individual $m \in M$ it is not true that for all $i \in M$, L(i, m).

3 marks

Exercise 2. Apply the semantic tableaux procedure, using the *breadth-first* search strategy, to show the validity of the following statement:

"If someone is loved by everybody, then everybody loves someone"

3 marks

Hint: We may use the *antecedent* Γ and the *succedent* Δ of a sequent $\Gamma \Rightarrow \Delta$ to represent a *queue*: we invert the rule which applies to the leftmost formula in the antecedent and in the succedent (to the left and to the right of the \Rightarrow sign).¹ To implement the *breadth-first* [or the *depth-first*] strategy it suffices to add the new formulas to the end of [in front of] the queue in the antecedent and in the succedent. Here are the rules for the quantifiers in the *breadth-first* strategy and then the rules for the universal quantifier in the *depth-first* strategy:

¹Here we assume that in a sequent $\Gamma \Rightarrow \Delta$, the antecedent Γ and the succedent Δ are lists of formulas. The letters a_0, \ldots, a_i to the left of each sequent are names of generic individuals declared so far, with the aim of making all formulas in Γ true and all formulas in Δ false.

Breadth-First search-strategy

$$\begin{array}{c} \underline{a_0, \dots, a_i, a_{i+1}; \Gamma \Rightarrow \Delta, A(a_{i+1})}_{a_0, \dots, a_i; \Gamma \Rightarrow \forall x.A, \Delta} \forall -\mathbf{R} \\ \hline \underline{a_0, \dots, a_i; \Gamma \Rightarrow \Delta, A(a_0), \dots, A(a_i), \exists x.A}_{a_0, \dots, a_i; (T, A(a_0), \dots, A(a_i), \forall x.A) \Rightarrow \Delta} \forall -\mathbf{L} \\ \hline \underline{a_0, \dots, a_i; \Gamma \Rightarrow \Delta, A(a_0), \dots, A(a_i), \exists x.A}_{a_0, \dots, a_i; (T, A(a_{i+1})) \Rightarrow \Delta} \exists -\mathbf{R} \\ \hline \underline{a_0, \dots, a_i; \exists x.A, \Gamma \Rightarrow \Delta} \\ \hline \underline{a_0, \dots, a_i; \exists x.A, \Gamma \Rightarrow \Delta} \exists -\mathbf{L} \end{array}$$

Depth-First search-strategy

$$\frac{a_0, \dots, a_i, a_{i+1}; \Gamma \Rightarrow A(a_{i+1}), \Delta}{a_0, \dots, a_i; \Gamma \Rightarrow \forall x.A, \Delta} \forall - \mathbf{R} \quad \frac{a_0, \dots, a_i; A(a_0), \dots, A(a_i), \forall x.A, \Gamma \Rightarrow \Delta}{a_0, \dots, a_i; \forall x.A, \Gamma \Rightarrow \Delta} \forall - \mathbf{L}$$

Exercise 3. (i) Apply the semantic tableaux procedure to the sequent

$$\Rightarrow (\exists z \forall v. B(v, z)) \lor (A \to A)$$

Use a *breadth-first* search procedure, starting from the leftmost formula.

2 marks

(ii) What happens if you use a *depth-first* search procedure, starting from the leftmost formula?

2 marks

TOTAL 10 MARKS